DIGITAL SIGNAL PROCESSING

SUBJECT CODE-BEE604

YEAR-III

SEMESTER –VI

INTRODUCTION TO DIGITAL SIGNAL PROCESSING WHAT IS DSP ? date back, saying as, water askin DSP is the processing of signals by digital . unadgelet but date vt means. In general terms, a signal is a stream of information representing anything from stock prices to data from a remote sensing satellite. In many cases, the signal is initially in the form of an analog electrical voltage or current, Produced for example by a microphone or some other type of transduces. In some situations the data is already in degotal form. eg: ofp from the readout system of a CD Allreagh the water interest the proof Player. du analog signal nust be converted into digital form before OSP techniques can be applied. An analog electrical vottage signal can be digitized using an analog to Digital converter. son analog segnal on sampling results in a descrete signal followed by quantization and encoding in order to convert the discrete signal to digital signal signals need to be processed in a variety of ways. For example, the output signal from a bransduces neary well be contaminated with noise. Processing the signal using a felter corcurt can remove at ableast reduce the unwanted part of the spipnel.

DSP technology is commonly employed nowadays in derices such as mobile phones, multimedia computers, Video seconders, CD players, hard disk drive controllars and moderns, and well soon replace analog circuitry TV sets and telephones.

An important application of DSP is in Signal Compression and decompression.

In CD Systems, for example, the music recorded on the CD is is compressed form (to increase storage capacity) and must be decompressed for the recorded signal to be reproduced.

Signal Compression is used in digital cellular phones to allows a greater no. of calls to be handled Simultaneously within each local "cell". Although the mathematical theory underlying DSP techniques Such as FFT, Masselet transform, Hilbert transform, Digital filter design and Signal Compression can be facely Complex, the memorical operations required to implement these techniques are in fait very Simple.

The architecture of DSP chip is designed to carry out such operations increditaly fast, processing upto tens of millions of samples par second, to provide real time performance.

In signal processing, the function of a filter is to remove unwanted parts of the signal, There are two main kinds of fulters, analog and digital. An analog feiter uses analog dechronce councils made from components such as nexestors, capacitors and op-amps to produce the required fulturing effect Such filter criecuits are widely used in applications such as noise reduction, video signal enhancement, graph equalizers in this filter and may other accos.

A digital filter uses a digetal processor to perform memorical calculations on Sampled values of the Signal. The processor may be a general purpose computer such as PC, a a Specialized DSP chip.

The main advantages of depotal follows over analog

1. A dégital filter is programmeable

R. Digital filters are easily designed, tested and implemented on a general purpose computer à workstation.

3. The characteristics of analog forthe concurts are Subject to darft and are dependent on temperature. Digital follows do not suffer from these problems, and so are extremely stable with respect to both tome and temperature.

A. Unloke their analog Computers Counterparts, digital filters can handle how frequency signals accurately. its the speed of DSP technology continues to increase digital filters are being applied to high frequency. Signals in the RF domain, which in the past was the exclusive preserve of analog technology.

5. Digital filters are very much more versatule in them ability to process signals in a variety of ways. is Digital filter adapt to changes in the characteristics of

UNIT-I

SIGNALS AND SYSTEMS

SIGNALS

A signal is defined as a function of one or more variables and Conveys information.

I signal is a physical quantity that varies with time in general, or any other independent Variable.

YONE - DIMENSIONAL SIGNAL

When a function depends on a single variable to represent the segnal, it is said to be 10 signal. Eg: ECG signal, Speech segnal.

117 TWO - DIMENSIONAL SIGNAL

When a function depends on two variables to represent the signal, it is said to be 2D signal. Eg: Video signal

111 MULTIDIMENSIONAL SIGNAL

When a function depends on more than two variables to represent the signal, it is said to be a multidemensional signal.

Eg: 3-Dimentional representation of images. SYSTEMS I system is defined as a physical derice that

that performs an operation on a signal.

I systero is an entity that manopulates one or more inport signals to perform a function, which results in a new output signal.



Processing of degital signals on a degital Computi is known as Digital Signal Processing. To perform the processing degetally, there is a need for an interface between the analog signal and the digital processor. This interface is called as an analog to digital (A/D) convertor. The opp of the A/D convertor is a degotal signal that is appropriate an an input to the degotal processor

The Digital Signal processor may be a digital Computer or a Small inveroprocessor programmed to Prifain the descend operations on the input signal. Programmable machines provide the flexiboloty to change the signal processing operations through a change in the software, whereas hardwired machines are difficult to reconfigure.

On the other hand, when signal processing operations are well defined, a hardwrited implementation which are more convernment. releable and runs faster than its programmable counterpart.

The analog to degital to analog (D/A) convertor is used to provide analog inform signal from the digital output of the Digital Signal processor. Example of analog ofp is speech signal (vorce) ADVANTAGES OF DIGITAL SIGNAL PROCESSING

*Flexibility in reconfiguring the digital signal processing operations. simply by changing the program. *Provides much better control of accuracy requirements.

* Digital Signals are easily stored in a magnetic media without loss of signal fideloty beyond that introduced in the A/D conversion!

* cheaper

APPLICATIONS

* Tale communications (Hodems, Veideo conferencing, etc. * Consumer Electronics (Digital Audro/TV, sound record * Instrumentation and Control. (Digital Fitter, PLL) * Image processing (Image Compression, enhancement) * Hedecine (X-ray Scanning, spectrum analysis of ECG and EEG signals to detect. the various disorders in heart & brain) * Speech processing

Sevenio logy (Detection of underground nuclear explosion & eastiquake monitoring) * Merletory

about at been in

LIMITATIONS

* System complexity

*Bandwidth limited by Sampling rate

* pouse consumption

CONCEPTS OF FREQUENCY IN ANALOG AND DIGITAL

to the concept of time. vovace, T = 1 is the French of the formation of the Se have but Hence both for continuous and discide time signals the nature of time(t) offeets the nature of frequency (f). ANALOG SIGNAL A simple harmonic oscollation is ives the more mathiniatically descerbed by the following continuous time semusoodal segual x(t)=Acos (at+0), -oxtxa A - Anylitude of the senusoed N-Frequency in radfeed. (N=arf) 0 - Phase in radians. f-Frequency in cycles per Second (a) Hz In-terms of F, acts can be written as estation as estate re(t)=vacos(arpt+0), -wataw $\frac{1}{2} = \frac{1}{2} = \frac{1}$ Acoso For valueration Sad sou Nageline and t 0 11 carge for analog 2mar ANALOG SINUSOIDAL SIGNAL

PROPERTIES DOLANA NU VONSURE 99 30 20930 HOD * For every fixed value of frequency F, x 12 is periodice . The is preserved to have ant x(t+T) = x(t) and a transmit where, T = 1 is the fundamental period of the since -sordal Signal. anternous and description * Continuous - time simusoidal signals with distinct frequencies are thenselves destanct (different) # Increasing the Frequency F results in an increase in the rate of oscellation of the signal. ives the more periods are included in a grown time interval. VF= 1/7. When F=0, T= ~ Due to continuity of the time variable t, we can increase the frequency F, without loved, with a corresponding increase in the rate of oscillabion. The sumuspidal signals carry over to the class of complexe exponential can be described as reto = A eicht +0) $\therefore x(t) = \sqrt{(\cos(t+0) + i)} \sin(t+0)$ $x(t) = A \cos(xt+\theta) = A e^{j(xt+\theta)} + A e^{-j(xt+\theta)}$ · · · coso= eltet For mathematical convenuence, we use both nogative and positive frequencies. Hence the fuquency range for analog sinusoed is ros x F x as.

SAMPLING Sampling is a process by which a continuous. signal is converted into a sequence of describe samples, with each sample representing the ampletude of the signal at a particular instant of time. Sampling is described by the relation $\sum_{n=1}^{\infty} \sum_{n=1}^{\infty} \sum_{n$ se (0) - descrete terre signal racto - Analog signal T-sampling period (or) Sampling interval Sampling, Fs = 1/7 · samples per second (09) To is the sampling Frequency Cherty) For a continuous time signal, the frequency Periodec sampling establishes a celationshop ofno the time raceables t & n of continuous-time and discrete time signals, respectively. t=nt= As a consequence of above equation, there excest a relationship blu the frequeny of analog Signal (2) + frequency of discrete signals. Consider an analog Scinusoidal signal, $x_{a}(t) = A \cos(axFt+0) \longrightarrow 0$



 $\frac{1}{2} = -\pi F_{S} \leq \mathcal{J} \leq \pi F_{S} = \frac{\pi}{2}$

. .

Periodic sampling of a continuous-time signal implies a mapping of the infonite frequency range for the variable F (or r) ento a finite frinct fuquency range for the variable f (or w)

Since the highest frequency is a descrete time signal is when a $f = \frac{1}{2}$, it follows that, with a Sampling rate Fs, the corresponding highest values of F and 2 are $F_{mon} = \frac{F_s}{2} = \frac{1}{27}$

If I'max = This = T If I'm is considered to be the maximum or highest fuguency in the input signal, these according to Nyquest, the Sampling rate Is should be greater than or equal to there the maximum frequency

FSZZFM. Talad

When Fs & 2 Fm., the signal cannot be reconstructed fully. Hence alwasing occurs.

This alrasing can be avoided if the input signal frequeneses are below one half of the Sampling frequency. This frequency Fig is called the Nyquest frequency for WN.

The impulse tracin of pulses can be expressed as

$$\begin{split} & \delta(t) = \sum_{n=0}^{\infty} \delta(t - nT_{0}) & \longrightarrow 0 \\ & \lambda_{0}(t) \ can be expressed watheristically as
$$& T_{0}(t) = \sum_{n=0}^{\infty} \delta(t - nT_{0}) \pi(nT_{0}) & \longrightarrow 0 \\ & \lambda_{0}(t) \ can be expressed watheristically as
$$& T_{0}(t) = \sum_{n=0}^{\infty} \delta(t - nT_{0}) \pi(nT_{0}) & \longrightarrow 0 \\ & \lambda_{0}(t) \ can be expressed sampled version of $\pi(t)$.
The Fourier Transform of impulse train of eqn 0 5
given as

$$& = \int_{0}^{\infty} \sum_{n=0}^{\infty} \delta(t - nf_{0}) & \longrightarrow \int_{0}^{\infty} \sum_{n=0}^{\infty} \sum_{n=0}^{\infty} \sum_{n=0}^{\infty} \frac{1}{T_{0}} \sum_{n=0}^{\infty} \sum_{n=0}^{\infty} \frac{1}{T_{0}} \sum\underbrace{1}{T_{0}} \sum\underbrace{1}{T_{0}}$$$$$$$$

SAMPLING THEOREM Let sails is a band limited signal with Xa(Jr): for INIXAM. Theor rails is unequely determined from its samples x(n)=xa(nT) if the sampling frequency Fs >2 Fm i.es sampling frequency must be at least the highest frequency present in the signal. PROOF. Consider alts as a input analog Signal. It has finite energy and finite duration. This xits is band limited signal. TX(E) Hore a for & the formation Transfor The above equipment shirt the last sampling a constructure fine p CONTINUOUS TIME SIGNAL W(E) 8(t-nTs) aprile a compe A MANA MANA (1) x = (1- 2) x <0.7 0 12 27 36 47 52 A UNIT INPULSE TRAIN USED AS A SAMPLING PUNCTION heat IX at the lat x at the tox at + cox at = cos x The above expression started that the game & the same graphing at - ale manage fieldand Eqn (3) and also the weet from

SAMPLED VERSION OF SIGNAL X(E).

$$\begin{array}{l} \therefore \times \langle f \rangle = \frac{1}{\partial M} \prod_{n=\infty}^{\infty} \chi\left(\frac{n}{\partial N}\right) e^{\int 2\pi f n/M} \qquad : T_{E} = \frac{1}{\Delta H}, \\ \times \langle f \rangle = \frac{1}{\partial M} \prod_{n=\infty}^{\infty} \chi\left(\frac{n}{\partial N}\right) e^{\int \pi f n/M} \qquad : T_{E} = \frac{1}{\Delta H}, \\ \times \langle f \rangle = \frac{1}{\partial M} \prod_{n=\infty}^{\infty} \chi\left(\frac{n}{\partial M}\right) e^{\int \pi f n/M} \qquad f n + N \leq f \leq M \longrightarrow \mathcal{O} \\ \end{array}$$

$$\begin{array}{l} \chi(f) & is obtained from above equ \chi(f) by taking IFT. \\ \chi(f) = 1FT(\chi(f)) = 1FT\left(\frac{1}{\partial N} \prod_{n=\infty}^{\infty} \chi\left(\frac{n}{\partial N}\right) e^{\int 2\pi f n/M} \int_{X} \\ n & \\ \end{array}$$

$$\begin{array}{l} \chi(f) = \frac{1}{\partial M} \prod_{n=\infty}^{\infty} \chi\left(\frac{n}{\partial M}\right) e^{\int 2\pi f n/M} \int_{X} \\ \frac{1}{\partial N} \prod_{n=\infty}^{\infty} \chi\left(\frac{n}{\partial M}\right) e^{\int 2\pi f n/M} \int_{X} \\ \frac{1}{\partial N} \prod_{n=\infty}^{\infty} \chi\left(\frac{n}{\partial M}\right) e^{\int 2\pi f n/M} \int_{X} \\ \frac{1}{\partial M} \prod_{n=\infty}^{\infty} \chi\left(\frac{n}{\partial M}\right) e^{\int 2\pi f n/M} \int_{X} \\ \frac{1}{\partial M} \prod_{n=\infty}^{\infty} \chi\left(\frac{n}{\partial M}\right) e^{\int 2\pi f n/M} \int_{X} \\ \frac{1}{\partial M} \prod_{n=\infty}^{\infty} \chi\left(\frac{n}{\partial M}\right) e^{\int 2\pi f n/M} \\ \frac{1}{\partial M} \prod_{n=\infty}^{\infty} \chi\left(\frac{n}{\partial M}\right) \frac{1}{\partial M} \prod_{M} e^{\int 2\pi f n/M} \\ \frac{1}{\partial M} \prod_{n=\infty}^{\infty} \chi\left(\frac{n}{\partial M}\right) \frac{1}{\partial M} \prod_{M} e^{\int 2\pi f n/M} \\ \frac{1}{\partial M} \prod_{n=\infty}^{\infty} \chi\left(\frac{n}{\partial M}\right) \frac{1}{\partial M} \prod_{M} e^{\int 2\pi M f n/M} \\ \frac{1}{\partial M} \prod_{n=\infty}^{\infty} \chi\left(\frac{n}{\partial M}\right) \frac{1}{\partial M} \prod_{M} e^{\int 2\pi M f n/M} \\ \frac{1}{\partial M} \prod_{m=\infty}^{\infty} \chi\left(\frac{n}{\partial M}\right) \frac{1}{\partial M} \prod_{M} e^{\int 2\pi M f n/M} \\ \frac{1}{\partial M} \prod_{m=\infty}^{\infty} \chi\left(\frac{n}{\partial M}\right) \frac{1}{\partial M} \prod_{M} e^{\int 2\pi M f n/M} \\ \frac{1}{\partial M} \prod_{m=\infty}^{\infty} \chi\left(\frac{n}{\partial M}\right) \frac{1}{\partial M} \prod_{M} e^{\int 2\pi M f n/M} \\ \frac{1}{\partial M} \prod_{m=\infty}^{\infty} \chi\left(\frac{n}{\partial M}\right) \frac{1}{\partial M} \prod_{M} e^{\int 2\pi M f n/M} \\ \frac{1}{\partial M} \prod_{m=\infty}^{\infty} \chi\left(\frac{n}{\partial M}\right) \frac{1}{\partial M} \prod_{M} e^{\int 2\pi M f n/M} \\ \frac{1}{\partial M} \prod_{M=\infty}^{\infty} \chi\left(\frac{n}{\partial M}\right) \frac{1}{\partial M} \prod_{M=\infty}^{\infty} \chi\left(\frac{n}{\partial M}\right) \frac{1}{\partial M} \\ \frac{1}{\partial M} \prod_{M=\infty}^{\infty} \chi\left(\frac{n}{\partial M}\right) \frac{1}{\partial M} \prod_{M=\infty}^{\infty} \chi\left(\frac{n}{\partial M}\right) \frac{1}{\partial M}$$

20 epn (5, first term represents spectrum that would have been obtained without sampling and rest of the terms under Summation represents spectrums repeating multiple fuquencies of sampling dequency fo.



W.K.T, FT [r_{glb}] = x_{glf} .

X(f) = 1 = 2(0Ts) e-JentinTs. Where xs(f) = 2 x (0Ts) e int The above eqn gives F-T of xits to terms of x(0)?.)



RECONSTRUCTION OF self FROM ITS SAMPLES

consider the signal xalt, = 10 cos 2 x (1000)t + 5 cos 2x (5000t) is to be Sampled.

is Determine the Nyquest rate for thus signal. is SIG the signal is sampled at 4 KHZ, will the signal be recovered from its samples?

solution :

") The given signal contains two casine wave FI= LOOOHZ and AI= 10 F2 = 5000 H2 and A2 = 5

" FMax = 5000 H2 .

Nyquist rate Fs Zd Fmax.

Fs ≥ 2(5000)

it's If the signal is sampled at 4 KHZ, the stignal can not be recovered from its samples; because Fs >2Fman

iv sequence Representation
A first duration sequence with time paymined
indicated by the symbol
$$h$$
 is represented as
 $\pi(n) = \int (1, 2, 2, 0, 0.5, 1.5)^2$
 dn infinite duration sequence can be
represented as
 $\pi(n) = \int \dots p(5, 1, 1-2, 3, 5) \dots f^2$
 d finite duration sequence that
satisfies the condition $\pi(n) = 0$ for $n < 0$ can be
represented as
 $\pi(n) = \int \partial_1 A_1 \partial_1 B_1 - S \int \partial_1 A_2 \partial_2 B_1 + B_2 \partial_2 B_1 + B_2 \partial_1 B_2 - B_1 + B_2 - B_1 + B_1 + B_2 \partial_1 B_2 - B_1 + B_1 + B_2 + B_1 +$











Vi} ADDITION OPERATION

$$\frac{x_1(n)}{x_2(n)}$$
 $\xrightarrow{y_1(n)} + x_2(n)$
 $x_3(n)$
Eg: $x_1(n) = \int 1, 3, 2, 5 \cdot 3$ $\pi_2(n) = \int 3, 1, 2, 4 \cdot 3$
 $y(n) = x_1(n) + \pi_2(n) = \int 4, 4, 4, 9 \cdot 3$
DISCRETE THE SYSTEMS
d discrete time gister is a device of
an algorithm that operates on a discrete time
input signal $x(n)$, according to some well defound
such, to produce autother descrete time Signal $y(n)$
called the output Signal.
 $y(n) = T [\pi(n)]$
 $x(n)$ DISCRETE THE y(n)
 $x(n)$ DISCRETE THE systems
 $y = x_1 + 1 + x_1 +$

is causal and Noncausal systems d system is said to be causal if the ofp of the system at any time n depends only at present and past exputs, but does not depend on future exputs. y(n) = F[x(n), x(n-1), x(n-2), ..., J.

Eg: causal System

 $y(n) = \alpha(n) + \alpha(n-1)$

Non-causal System

anitaling y(n) = x(an) in and but with open an

iliy LINEAR AND NON-LINEAR SYSTEMS

I System that satisfies the superposition prenciple is said to be a tinear System.

superposition principle!

It states that the response of the system to a weighted sum of signal should be equal to the corresponding weighted sum of the outputs of the system to each of the individual signal.

 $T[a_1 x_1(n) + a_2 x_2(n)] = a_1 T[x_1(n)] + a_2 T[x_2(n)]$

iv) THE VARIENT AND THE INVARIENT SYSTEMS d system is said to be time invarient or shift invarient if the characteristics of the system do not change with time.

 $y(n) \leftrightarrow T[x(n)]$ $y(n-k) \leftrightarrow T[x(n-k)]$



If the two systems are connected in parallel,
then the ownell impulse response is qual to sum
of two impulse responses.
ii) CASCADE CONNECTION OF TWO SYSTEMS

$$x(n) = h_1(n) \frac{y(n)}{h_2(n)} \frac{y(n)}{y(n)} \frac{y(n)}{x(n)} \frac{y(n)}{h_1(n) * h_2(n)} \frac{y(n)}{y(n)}$$

 $y_1(k) = x(k) * h_1(k)$
 $= \underbrace{z}_{k=-\infty} x(k) h_1(n-k) + h_2(k)$
 $= \underbrace{z}_{k=-\infty} x(k) h_1(n-k) + h_2(n)$
 $= \underbrace{z}_{k=-\infty} x(k) h_1(n-k) + h_2(n)$
 $= \underbrace{z}_{k=-\infty} x(w) + h_2(k)$
 $= \underbrace{z}_{k=-\infty} x(w) h_1(k-w) + h_2(n)$
 $= \underbrace{z}_{k=-\infty} x(w) + h_2(n-w) + h_2(n-w) + p)$
 $= \underbrace{z}_{k=-\infty} x(w) = h_1(p) h_2(n-w+p)$
 $= \underbrace{z}_{k=-\infty} x(w) = h_1(p) h_2(n-w+p)$
 $= \underbrace{z}_{k=-\infty} x(w) = h_1(k) h_2(n-w+p)$
 $= h_1(m) * h_2(m)$.
Hence the impulse response of two it TI
Systems Connected in caseade is the convolution

of the individual impube response.

Determine the values of power and energy of the following
signals. Find whether the signals are power an energy signal.
i)
$$x(n) = (\frac{1}{2})^{n} u(n)$$
 ii $x(n) = d^{2} [\frac{n}{2}n + \frac{n}{2}n]$ iii $x(n) = sin (\frac{n}{2}p)$
solution:
i) Griven $x(n) = (\frac{1}{2})^{n} u(n)$
The energy of the signal
 $E = \sum_{n=0}^{\infty} 1 x(n)^{2}$
 $= \sum_{n=0}^{\infty} [\frac{n}{2}n]^{n}$
 $= \frac{1}{1-\frac{1}{2}} = \frac{9}{8}$
The power $p = \lim_{n \to 0} \frac{1}{n} \lim_{n \to 0} \frac{1}{2n} [\frac{1-(\frac{1}{2}n)^{n}}{1-\frac{1}{2}}]$
 $= 0$
The energy is first and power is zero.
 \therefore the signal is an energy signal.
W/ $x(n) = e^{2[\frac{n}{2}n+\frac{n}{2}n]}$
 $E = \sum_{n=0}^{\infty} [\frac{1}{2}(\frac{n}{2}n+\frac{n}{2}n]^{2}$
 $= 0$
The maximal is an energy signal.
 $W/ x(n) = e^{2[\frac{n}{2}n+\frac{n}{2}n]}$
 $E = \sum_{n=0}^{\infty} [e^{2[\frac{n}{2}n+\frac{n}{2}n]}]^{2}$
 $= \lim_{n \to 0} \frac{1}{2} \lim_$

iii)
$$\pi(m) = \sin(\pi)$$

 $E = \sum_{n=\infty}^{\infty} |\sin(2(\pi_{n}n))| = \sum_{n=\infty}^{\infty} \left[\frac{1-\cos(\pi_{n}n)}{2}\right] = \infty$
 $P = \lim_{N \to \infty} \frac{1}{2N+1} \sum_{n=n}^{\infty} |\sin(2(\pi_{n}n))|$
 $= \lim_{N \to \infty} \frac{1}{2N+1} \sum_{n=n}^{\infty} \frac{1-\cos(\pi_{n}n)}{2} = \frac{1}{2} \lim_{N \to \infty} \frac{1}{2N+1} \sum_{n=n}^{N} \frac{1}{n-n}$
 $= \frac{1}{2}$.
This signal is power signal.
Determine whether or not each of the following signals is
periodic. If a signal is periodice, specify its fundamental
Period.
 $(\frac{1}{2}\pi(n) = e^{i\beta\pi n})$ $\frac{1}{2}\pi(n) = e^{i\beta\pi(n+\frac{1}{2})}$ $\frac{1}{2}(\frac{1}{2}\pi(n) = \cos(\frac{1}{2}\pi)n$.
 $(\frac{1}{2}\pi(n) = e^{i\beta\pi n})$ $\frac{1}{2}\pi(n) = \cos(\pi)$ $\frac{1}{2}\pi(n) = \cos(\frac{1}{2}\pi)n$.
 $N = \pi(\pi)$ $\frac{1}{2}\pi(n) = \cos(\pi)$ $\frac{1}{2}\pi(n) = \cos(\pi)$.
Solution:
 $\frac{1}{2}\pi(n) = e^{i\beta\pi n}$ $\frac{1}{2}\pi(n) = \cos(\pi) + \cos(3\pi)n$.
Solution:
 $\frac{1}{2}\pi(n) = e^{i\beta\pi n}$ $\frac{1}{2}\pi(n) = \cos(\pi)$ $\frac{1}{2}\pi(n) = e^{j\omega n}$.
 $\frac{1}{2}\pi(n) = e^{i\beta\pi n}$ $\frac{1}{2}\pi(n) = \cos(\pi)$ $\frac{1}{2}\pi(n) = e^{j\omega n}$.
 $\frac{1}{2}\pi(n) = e^{i\beta\pi}(n+\frac{1}{2}) = 2\pi(\frac{m}{6\pi})$.
The number value of m for wheth N is integer is 3.
 $N = 2\pi(\frac{3}{2}\pi) = 1$.
 $\frac{1}{2}\pi(n) = e^{i\beta\pi}(n+\frac{1}{2})$ $\frac{1}{2}\pi(n) = e^{i\beta\pi}(n+\frac{1}{2})$
 $\omega_{0} = \frac{3}{2}\pi(n+\frac{1}{2})$ $\omega_{0} = 2\pi$.
 $\frac{1}{2}\pi(n) = \cos(2\pi)n$ $\frac{1}{2}\pi(n) = 2\pi$.
 $\frac{1}{$

Solution:
i's
$$y(n) = z(n) + \frac{1}{x(n-1)}$$

For $n=-1$; $y(n) = z(n) + \frac{1}{x(n)}$
For $n=0$; $y(0) = z(n) + \frac{1}{x(n)}$
For $n=1$; $y(n) = z(n) + \frac{1}{x(n)}$
For $n=1$; $y(n) = z(n) + \frac{1}{x(n)}$
For $n=1$; $y(n) = z(n)$
For $n=0$; $y(n) = z(n)$
For $n=0$; $y(n) = z(n)$
For $n=1$; $y(n) = z(n) + \frac{1}{x(n-1)}$
For two ilp sequences $z_1(n) = z_3(n)$ the corresponding of $y(n) = y(n) = z(n) + \frac{1}{x(n-1)}$
The of p due to weighted cause of inpute is
 $y_3(n) = T(z_1(n) + a_2 x_3(n)) + \frac{1}{a_1 x_1(n-1) + a_2 x_2(n-1)} \rightarrow 0$

The linear Combination of the two output is

$$a_1y_1(n) + a_2y_2(n) = a_1x_1(n) + \frac{a_1}{x_1(n-1)} + a_2x_2(n) + \frac{a_2}{x_2(n-1)}$$

 $equils (n = \pm (3)$
 \cdots The system is non-linear.
 $ii' > y_1(n) = T[e_1(n)] = x_1^2(n)$
 $y_1(n) = T[e_1(n)] = x_1^2(n)$
 $y_2(n) = T[e_1(n)] = x_1^2(n)$
 $y_2(n) = T[e_1(n)] = x_1^2(n)$
 $y_2(n) = T[a_1x_1(n) + a_2x_2(n)] \xrightarrow{a} = a_1x_2(n) + a_2x_2(n) \xrightarrow{a} = 0$
 $a_1 T[a_1x_1(n) + a_2x_2(n)] = [a_1x_1(n) + a_2x_2(n)] \xrightarrow{a} = 0$
 $a_1 T[a_1(n)] + a_2x_2(n)] = a_1x_1^2(n) + a_2x_2(n) \xrightarrow{a} = 0$.
 $\leq qn's = 0 \pm e$
 \therefore The system is non-linear.
 $iii' > y = 0 = n x_1(n)$
 $y_2(n) = T[a_1(n)] = n x_2(n)$.
 $y_3(n) = T[a_1(n)] = n x_2(n)$.
 $y_3(n) = T[a_1(n)] = n x_2(n)$.
 $y_4(n) = T[a_1(n)] + a_2x_3(n)] = a_1n x_1 = a_2x_2(n) \longrightarrow 0$.
 $a_1 T[x_1(n)] + a_2x_3(n)] = a_1n x_1 = a_2x_2(n) \longrightarrow 0$.
 $e = 0$, \therefore The system is linear.
Determine if the following systems are time-convaluent or
time valuent
 $e^{y} y_{(n)} = x_{(n)} + x_{(n-1)}$.
 $i' Given y_{(n)} = x_{(n)} + x_{(n-k-1)}$.
 $i' Given y_{(n)} = x_{(n)} + x_{(n-k-1)}$.
 $i' we delay the of ey k units in time then
 $y_{(n,k)} = T[x_{(n-k)}] = x_{(n-k)} + x_{(n-k-1)}$.$

6}
UNIT- 2

Z TRANSFORM of linear-time invarient describe time systems in the grequency domain? dr (2), x 0 = for a dread, co In z-domain the convolution of two time domain signals is equilant to multiplication of their Corresponding z- transform. DEFINITION The z-transform of a descrete time sognal sens is defined as X(Z)= Z x(n) X (Two ended X-transform). where x is a Complexe variable In polar form & can be expressed as Z = rein conde = [con-obe] = where & is the radius of the cilecte. × (rein) = = r(n) rhein reduces to For r=1, the above expression Fourier transform of xins. res the x-transform evaluated on the unit circle corresponds to the Fourier transform. If sichs is a causal sequence i'es secoso for n=0, then the x-transform is in X,(x) = 2 x(n) x (One sided z-transform)

PROPERTIES OF THE X-TEAMSPORM
i) LINEARITY
If
$$x_1(x) = x\{x_1(n)\}$$
 and $x_2(x) = x\{x_3(n)\}$, then
 $x\{a x_1(n)+b x_2(n)\} = a x_1(x)+b x_2(x)$
PROOF:
 $x\{a x_1(n)+b x_2(n)\} = \sum_{n=-\infty}^{\infty} [a x_1(n)+b x_2(n)] x^{n}$
 $= a \sum_{n=-\infty}^{\infty} x_1(n) x^{n} + b \sum_{n=-\infty}^{\infty} x_2(n) x^{n}$
 $= a x_1(x) + b x_2(x)$
ii) TIME SHIFT OR TRANSLATION
Sf $x(x) = x\{x(n)\}$ and the initial conductions
for $x(n)$ are xeros. Then
 $z\{x(n-m)\} = x^{n} x(x)$
FROBF:
 $x\{x(n-m)\} = x^{n} x(n) x^{n}$
Let $(n-m) = 1$, then
 $x\{x(n)=x\}\} = x^{n} x(x)$
iii) HOLTIPLICATION BY AN EXPONENTIAL SERVENCE
 $g\{x(n)=x\{x(n)\}\} = x^{(n-1)} x^{(n)}$
 $x\{a^n x(n)\} = x(a^{(n)}x)$
 $x\{a^n x(n)\} = x(a^{(n)}x)$

Where Reversate
If
$$\chi(x) = z\{\pi(m)\}$$
, then
 $\chi\{\pi(-n)\} = \chi(z^{-1})$
Proof:
 $\chi\{\pi(-n)\} = \sum_{n=-\infty}^{\infty} \chi(n) z^{-n}$, where $l=-n$
 $= \sum_{n=-\infty}^{\infty} \chi(n) (z^{-1})^{-1}$, where $l=-n$
 $l=-\infty$
 $= \frac{z}{\lambda} = \chi(z^{-1})$
where the ROC is $k_{2} \le lz \le \frac{1}{k_{1}}$
 Y DIFFEEENTIATION OF $\chi(x)$
 $Q = \chi(x) = \chi\{\pi(m)\}$, then
 $-\chi\{\pi,\chi(n)\} = -\chi \frac{d}{d\chi}\chi(x)$
PROOF:
 $\chi(x) = \sum_{n=-\infty}^{\infty} \chi(n) \frac{d}{d\chi} x^{-n}$
 $d_{\chi\chi}(x) = \sum_{n=-\infty}^{\infty} \chi(n) \frac{d}{d\chi} x^{-n}$
 $= -\frac{1}{\chi} \sum_{n=-\infty}^{\infty} \chi(n) \chi^{-n}$
 $= -\frac{1}{\chi} \sum_{n=-\infty}^{\infty} \chi(n) \chi^{-n}$
 $-\chi \frac{d}{d\chi}\chi(x) = \sum_{n=-\infty}^{\infty} n\chi(n) \chi^{-n}$
 $-\chi \frac{d}{d\chi}\chi(x) = \sum_{n=-\infty}^{\infty} n\chi(n) \chi^{-n}$
 $= -\chi\{\pi,\chi(n)\}$

VIL CONVOLUTION THEOREM SLITT WI If x(x)=x{a(n)}, and H(x)= 25 h(n)}, then $\chi f x(n) * h(n) f = x(x) H(x) [(n) + x)$ n(n) * h(n) -> Linear Convolution Sequences. PROOF: We have y(n) = 2(n) + h(n) and $= \underbrace{\mathcal{S}}_{k=-\infty} \mathfrak{n}(k) h(n-k)$ $\gamma(x) = x \{y(n)\} = \underbrace{z}_{n=-\omega} \begin{bmatrix} z^{\omega} & x(k)h(n-k) \end{bmatrix} z^{-h}$ $= \underbrace{\underbrace{\overset{o}}{\underset{n=-\infty}{\overset{o}}}_{k=-\infty} \underbrace{\overset{o}}{\underset{n(k)}{\overset{\tau}{\underset{k}}}_{k(k)} \underbrace{\overset{f}{\underset{k}}{\overset{k}{\underset{k}}}_{k(n-k)} \underbrace{\overset{f}{\underset{k}}{\underset{n(n-k)}{\overset{r}{\underset{k}}}}_{k(n-k)} \underbrace{\overset{f}{\underset{k}}{\overset{f}{\underset{k}}}_{k(n-k)} \underbrace{\overset{f}{\underset{k}}}_{k(n-k)} \underbrace{\overset{f}{\underset{k}}{\overset{f}{\underset{k}}}_{k(n-k)} \underbrace{\overset{f}{\underset{k}}}_{k(n-k)} \underbrace{\overset{f}{\underset{k}}}_{k$ Interchange the order of the summation Y(x)= = x(k) = k = h(n-k) x -(n-k) k= x n= n= = ~ x(k) z k ~ h(l) z l where l=n-k z { u(n) * h(n) } = x(z) H(z) (n)x = cx)x h VIIL PARSEVAL'S RELATION Let us consider two complex sequences 14(n) and 12(n). Parseval's relation states that $\sum_{n=-\infty}^{\infty} x_{1}(n) x_{2}^{*}(n) = \frac{1}{2\pi j} \oint X_{1}(u) X_{2}^{*}(\frac{1}{u^{*}}) u^{-1} du$ Where the contour of integration must be the overlap of the seguons of convergence of in $X_1(v)$ and $X_2^{*}\left(\frac{1}{v^{*}}\right)$

PROOF!
Let
$$y(n) = x_1(n) \times x^{m}(n)$$
. Then $y(n) = \sum_{n=-\infty}^{\infty} x_1(n) x_{n}^{m}(n) \times x^{m}(n)$
Will PARSEVAL'S RELATION
Parseval's relation states that the total average
power in a discret periodic signal $x(n)$, equals the sum
of the average period x signal $x(n)$, equals the sum
of the average period x^{m} is diverdiable hormonic.
Components, which is turn equals to squared meganitude of $x(n)$.
 $\sum_{n=-\infty}^{\infty} |x(n)|^{2} \ll T \to \frac{1}{2\pi i} \int |x(n)|^{2}$
PREOF
If $x(n)$ is the discret periodic signal is given by
 $\sum_{n=-\infty}^{\infty} |x(n)|^{2} = \sum_{n=-\infty}^{\infty} x(n) [x(n)]^{n}$
 $\sum_{n=-\infty}^{\infty} |x(n)|^{2} = \sum_{n=-\infty}^{\infty} x(n) [x(n)]^{n}$
Differencianging the position of integral and summation,
 $\sum_{n=-\infty}^{\infty} |x(n)|^{2} = \sum_{n=-\infty}^{\infty} x^{n}(x) \times (x)$
 $= \frac{1}{2\pi i} \int |x(x)|^{2}$
 $= \frac{1}{2\pi i} \int |x(x)|^{2}$
 $\sum_{n=-\infty}^{\infty} |x(n)|^{2} = \sum_{n=-\infty}^{\infty} x^{n}(x) \times (x)$
 $= \frac{1}{2\pi i} \int |x(n)|^{2}$
 $\sum_{n=-\infty}^{\infty} |x(n)|^{2} = \sum_{n=-\infty}^{\infty} x^{n}(x) \times (x)$
 $= \frac{1}{2\pi i} \int |x(n)|^{2}$
Willy initial value theorem
 $x(n) = \sum_{n=-\infty}^{\infty} \sum_{n=-\infty}^{\infty} |x(n)|^{2}$

Z-TRANSFORM AND ROC OF FINITE DURATION SEQUENCES. The region of convergence (Roc) of x(x) is the set of all values of x for which x(x) attains a finite value. RIGHT HAND SEQUENCE

A right hand stoke sequence is one for which 2017=0 for all n < no where no is + ve or -ve but finite. If no is greater than or equal to zero, the resulting sequence is causal a a positive time sequence. For such type of sequence the roc is antike X plane except 200.

Eg:1

Find the z-transform and ROC of the causal sequence

$$\chi(0) = \{1, 0, 3, -1, 2\}$$

Solutión:

For causal Sequence res x(n)=0 for nx0, then the x transform is

$$(z) = \sum_{n=0}^{\infty} \chi(n) z^{n}$$

= $\chi(0) z^{-1} + \chi(1) z^{-1} + \chi(2) z^{-2} + \chi(3) z^{-3} + \chi(4) z^{-4}$

 $x_{+}(x) = 1 + 3x^{-2} - x^{-3} + 2x^{-4}$

The X+12) converges for all values of z except at z=0.

LEFT HAND SEQUENCE

2n=0 dos all n≥no

If no <0 the resulting sequence is anticausal Sequence.

For such type of sequence the ROC is entrice z-plane except at z= ~

Eg': 2 Find the z-transform and Roc of the autreausal Sequence. $207 = \{-3, -2, -1, 0, 1\}$ Solution. $x(x) = \frac{x}{2} x(x) x^{n}$ $\lambda(x) = 1 - x^2 - 2x^3 - 3x^4$ The x(x) converges for all values of x exceptal 2:00 TWO SIDED SEQUENCE. of segnal that has finite aluxation on both the lef and right hand seales is known as two scaled sequence For such type of sequence the ROC is entire x-plane except at x=0 and x=00 E9:3 Find the z-transform of the sequence 2 (1) = { 2, -1, 3, 2, 1, 0, 2, 3, -1 } Solution : $X(x) = \sum_{n=1}^{\infty} \pi(n) x^{-n}$ -2x4-z3+3x2+22+1+2x2+3x2-z4. The XX7 converges for all values of z except X=0 & X=0. Z-TRANSFORM AND ROC OF INFINITE DURATION SEQUENCE. Determine the X-transform and Roc of the signal $\chi(n) = \alpha^n \mu(n)$.

Solution
The given signal is causal and of firsts duration.
The x-transform of xim is given by

$$X(x) = \sum_{n=0}^{\infty} n \sin x^{-n}$$

 $= \sum_{n=0}^{\infty} a^n x^{-n} = \sum_{n=0}^{\infty} (ax^{-1})^n$
 $= \sum_{n=0}^{\infty} a^n x^{-n} = \sum_{n=0}^{\infty} (ax^{-1})^n$
 $\lim_{n \to 0} x^{-1} = 1 + x + x^2 + x^3 + \dots = \frac{1}{1-x}$ if $|z| | |z| | |x|$
 $\therefore x(x) = \frac{1}{1-ax^{-1}} = \frac{x}{x-a}$ Roc is $|z| > a$.
Les the Roc is the extension of a curcle thaving radius (a)
Eg: 5
Find the x-transform and the Roc of the signal
 $x(n) = -b^n u(-n-1)$.
Solution
The given signal is of infinite duration and
anticaused.
 $X(x) = \frac{x}{1-b} x^{-n} = -\frac{x}{b-a} (b^{-1}x)^n$.
The above searces converges to $b^{-1}x| < 1$,
 $x(x) = -\left[\sum_{n=0}^{\infty} b^{-1}x^{-n} - \int_{x-1}^{\infty} (b^{-1}x)^{-n}$.
The above searces converges to $b^{-1}x| < 1$,
 $X(x) = -\left[\sum_{n=0}^{\infty} b^{-1}x^{-n} - \int_{x-1}^{\infty} (b^{-1}x)^{-n}$.
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 $X(x) = -\left[\sum_{n=0}^{\infty} b^{-1}x^{-n} - \int_{x-1}^{\infty} (b^{-1}x)^{-n}$.
The above searces converges to $b^{-1}x| < 1$,
 $X(x) = -\left[\sum_{n=0}^{\infty} x - b^{-1}x^{-1} + b^{-1}x^{-1}x^{-1}\right]$
 $= \frac{x}{x-b}$ Roc: $|x| < b$.
The Roc is new the interior of a curcle.

Eg: 6 Fond the stability of the System whose impube response h(n) = 2ⁿ u(n) Solution: Griven h(n) = 2ⁿ u(n) Roc. Roc.

$$H(x) = \frac{x}{x-2} \quad |x| > 2$$

The ROC is 121>2. 14 does not contain const cucle.

PROPERTIES OF ROC

1. The ROC is a zing on dusc in the z-plane centural at the vergen.

2. The ROC cannot contain any poles

8. If 200 is a causal sequence the ROC is the entire Z-plane except at 200

2-plane except at 2=x.

5. If allow is a finite duration, two scaled sequence the ROC is contine 2 plane except at 200 \$ 200.

6. If surs is an enfinite duration, two scoled sequence the ROC will Contain consist of a ring in the z-plane, bounded on the intervor and exterior by pole, not containing any poles.

T. The Roc of a LTI stable System contains the must curcle.

8. The ROC nust be a connected region.

Find x-transform of the segnal
$$2(n) = [\frac{1}{2}(3)^n - 4(2)^n] u(n)$$
.
solution:

$$x(x) = \sum_{n=0}^{\infty} x(n) x^{-n}$$

$$= \sum_{n=0}^{\infty} [\frac{3}{2}(3)^n - 4(2)^n] u(n) x^{-n}$$

$$= \frac{3}{n=0} [\frac{3}{2}(3)^n - 4(2)^n] x^{-n}$$

$$= \frac{3}{n=0} (\frac{3}{2}(3)^n - 4\frac{3}{2n})^n (2x^{-1})^n$$

$$= \frac{3}{n=0} (\frac{3}{2}x^{-1})^n - 4\frac{3}{2n-0} (2x^{-1})^n$$

$$= \frac{3}{1-3x^2} - \frac{4}{1-3x^2} \quad \text{Roc}(1/x) > 3$$

$$= \frac{3x}{x-3} - \frac{4x}{x-2}$$

$$= \frac{3x}{x-3} - \frac{4x}{x-2}$$

$$= \frac{3x}{x-3} - \frac{4x}{x-2}$$
by Fund the z-transform of the sequence $x(n) = (\frac{1}{3})^{n-1}u(n-2)$
solution:

$$x[[\frac{1}{3})^n u(n)] = \frac{x}{x-y_3}$$
Using the shifting property

$$x[\frac{1}{2}(n-1)] = x^{-1}x(x)$$

$$\int f^{-1}u(n-1) = x^{-1}\frac{x}{x-y_3} = \frac{1}{x-y_3}$$

$$\sum Fund the z-transform of the sequence $x(n) = na^n u(n)$
solution:

$$x[\frac{n}{2}nun] = \frac{1}{1-ax^2} \quad |x| > |a|$$

$$\sum colution = na^n u(n)$$

$$x[n xn] = -x \frac{1}{dx} \times (x)$$

$$x[n xn] = -x \frac{1}{dx} \times (x)$$$$

$$= \frac{\alpha z^{-1}}{(1 - \alpha z^{-1})^2} : [z] > |a|$$

$$: -z \frac{d}{dz} (\frac{1}{1 - \alpha z^{-1}}) = z \left[\frac{1}{(1 - \alpha z^{-1})^2} + \alpha \frac{1}{z^2} \right]$$

$$= \frac{\alpha z^{-1}}{(1 - \alpha z^{-1})^2}.$$
4) Find inverse z-t constant of $x(z) = \log (1 - 0.5 z^{-1}); |z| > 0.5$
using Differentiation property.
Solution:
Gaiven $x(z') = \log (1 - 0.5 z^{-1})$
Differentiate on both Andes we get

$$\frac{d}{dz} x(z) = \frac{0.5 z^{-2}}{1 - 0.5 z^{-1}} \qquad \frac{d}{d\pi} \cdot 0 = \frac{1}{\pi}.$$
Multiply both sides: by $-z$ we obtain

$$-z \frac{d}{dz} x(z) = \frac{-0.5 z^{-1}}{1 - 0.5 z^{-1}}$$

$$Z[h z \ln^{2}] = -z \frac{d}{dz} x(z') = \frac{-0.5 z^{-1}}{1 - 0.5 z^{-1}}$$

$$= -0.5 z [(0 \cdot 5)^{h-1} u(n-1)] = \frac{1}{1 - 0.5 z^{-1}}$$

$$X[h z \ln^{2}] = -0.5 (0 \cdot 5)^{h-1} u(n-1)] = \frac{z^{-1}}{1 - 0.5 z^{-1}} = \frac{1}{1 - 0.5 z^{-1}}$$

5) Find the System function and impube response of the hydro
described by the difference equation
$$y(n) = \frac{1}{2}y(n-1) + 2(n)$$

solution:
(Given $y(n) = \frac{1}{2}y(n-1) + \frac{1}{2}(n)$
Taking x -transform on both sides we get
 $y(x) = \frac{1}{2}x^{-1}y(x) + x(x)$
 $y(x) - \frac{1}{2}x^{-1}y(x) = x(x)$
Now the System function is $H(x) = \frac{y(x)}{x(x)} = \frac{1}{1-\frac{1}{2}x^{-1}}$
By taking inverse x -transform
 $h(n) = (\frac{1}{2})^n u(n)$.
6) Find the System function and impulse response of
the system described by the difference equation
 $y(n) = x(n) + \frac{1}{2}x(n-1) - \frac{1}{2}x(n-3)$
Solution
Griven $y(n) = x(n) + \frac{1}{2}x(n-1) - \frac{1}{2}x(n-3) + \frac{1}{2}x(n-3)$
Paking x -transform on both sides
 $y(x) = x(x) + \frac{1}{2}x(x) - \frac{1}{2}x(x) + \frac{1}{2}x(x)$
 $y(x) = (1 + \frac{1}{2}x^{-1} - \frac{1}{2}x^{-\frac{3}{2}} + \frac{1}{2}x(x)$
 $y(x) = (1 + \frac{1}{2}x^{-1} - \frac{1}{2}x^{-\frac{3}{2}} + \frac{1}{2}x(x)$
 $y(x) = (1 + \frac{1}{2}x^{-1} - \frac{1}{2}x^{-\frac{3}{2}} + \frac{1}{2}x(x)$
 $h(n) = 1$, $h(n) = 2$, $h(n) = -4$, $h(n) = 1$.
 \therefore Despalse response, $h(n) = \int_{1}^{1} \frac{1}{2}(-4, -1)^{\frac{3}{2}}$

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Now the System function is $H(x) = \frac{y(x)}{x(x)} = \frac{1}{1-\frac{1}{2}x^{-1}}$
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 $h(n) = (\frac{1}{2})^n u(n)$.
6) Find the System function and impulse response of
the system described by the difference equation
 $y(n) = x(n) + \frac{1}{2}x(n-1) - \frac{1}{2}x(n-3)$
Solution
Griven $y(n) = x(n) + \frac{1}{2}x(n-1) - \frac{1}{2}x(n-3) + \frac{1}{2}x(n-3)$
Paking x -transform on both sides
 $y(x) = x(x) + \frac{1}{2}x(x) - \frac{1}{2}x(x) + \frac{1}{2}x(x)$
 $y(x) = (1 + \frac{1}{2}x^{-1} - \frac{1}{2}x^{-\frac{3}{2}} + \frac{1}{2}x(x)$
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$$\frac{x}{12} + \frac{7x^2}{14} + \frac{87}{1738}x^3$$

$$\frac{x}{12} - \frac{7x^4}{14} + \frac{7x^2}{1738}x^3$$

$$\frac{x}{12} - \frac{7x^4}{12} + \frac{87}{12}x^3$$

$$\frac{7x^4}{12} - \frac{7x^4}{12} + \frac{7x^3}{12}x^5$$

$$\frac{7x^4}{12} - \frac{7x^4}{12}x^4 + \frac{7x^3}{144} + \frac{7x^4}{144}$$

$$\frac{37x^3}{144} - \frac{7x^4}{1728}x^7 + \frac{87}{1728}x^5$$

$$x(x) = \frac{1}{12}x + \frac{7}{144}x^2 + \frac{37}{1728}x^3 + \dots$$

$$= \frac{31}{144} - \frac{1728}{1728}x^3 +$$

no find C₁
Sub
$$x = -1$$
, in april G
 $G = a$.
no find C₂
Sub $x = -2$ in eqn(a)
 $C_2 = -1$.
 $\therefore x(x) = \frac{3}{x_{+1}} - \frac{1}{x_{+2}}$.
 $x(x) = 2 \frac{x}{x_{+1}} - \frac{x}{x_{+2}}$.
 $ds \ Roc \ is \ |x| > 2 \ the sequence is causal and by$
interpretion, we
 $n(n) = a \in 1^n u(n) - (-a)^n u(n)$.
by Given $x(z) = \frac{x(x^2 - 4x + s)}{(z - 1)(z - 2)(z - 3)}$.
 $\frac{x(x)}{x} = \frac{x^2 - 4x + s}{(z - 1)(z - 2)(z - 3)}$.
 $\frac{x(x)}{x} = \frac{x^2 - 4x + s}{(z - 1)(z - 2)(z - 3)}$.
 $\frac{x(x)}{x} = \frac{x^2 - 4x + s}{(z - 1)(z - 3) + C_2(x - 1)(x - 3)} + C_3(x - 1)(x - 3)$.
No final c₁
Sub $x = 1$ en eqn (a)
 $C_1 = 1$.
No final c₂
Sub $x = 2$ in eqn (b)
 $C_2 = -1$.
No final c₃
Sub $x = 3$ in eqn (b)
 $C_3 = 1$.
 $\sum \frac{x(x)}{x} = \frac{1}{x - 1} + \frac{-1}{x - 2} + \frac{1}{x - 3}$.
 $x(x) = \frac{x}{x} - 1$.

1) In case when the ROC is 2<21×3, the signal 2007 is two sided. The poles x=1 and x=2 provide the causal part and the pole x=3 provides anticausal part.

.: 2(0)=u(0)-(3)ⁿu(0)-(3)ⁿu(-n-1)



AIm R?

Im(Z)

2-plane.

D

2-plany Roc

>Re(Z)

ii) In case when ROC is 1×1>3, the signal secons is causal and all three terms are causal terms.

いえのつこしのつ しょうれしの + (3)れしい



·ix(0) =- u (=n-1) + (2)ⁿ u (−n-1) - (3)ⁿ u (−n-1)

CONVOLUTION

If the input to the system is a unit impube i'es $\mathcal{R}(n) = \mathcal{S}(n)$, then impube response is

h(m) = T[s(m)]

For a linear time-invacient System. if the input sequence 2000 and impube response him are given, we can find the output y(n) by using the equation

$$y(n) = \overset{\sim}{\underset{k=-\infty}{\overset{\sim}}} x(k)h(n-k)$$

which is known as convolution sum and can be represented as

y(n)= x(n) * h(n)

PROPERTIES OF CONVOLUTION

is Commutative Law

2(n) * h(n) = h(n) * 2(n)

is Associative Law

iii) Distributive Law

$$x(n) + [h_1(n) + h_2(n)] = x(n) + h_1(n) + x(n) + h_2(n)$$

ii) CIRCULAR CONVOLUTION

Yens = xeins @ 22(n)

I If 2(n) contains & samples & 22(n) contains H samples, then the no. of samples in y(n) is the max (H,L)

$$\begin{array}{c} p \text{ beforming the convolution sums of two sequences} \\ x(n) = \int s_{1} s_{1} r_{1} r_{2} f_{1} f_{1}(n) = \int r_{1} s_{1} r_{2} f_{1} f_{2} f_{2}$$

$$\begin{array}{l} n_{21}, \\ y_{(1)} = \sum_{k=\infty}^{\infty} x_{(k)} h_{(1-k)} \quad k_{2} + 1/\rho_{1}/\rho_{1}^{3}, \\ h_{(1-k)} \\ = b_{(2)+\beta} (1) + 2(2) + 1(2) \\ = 8. \\ n_{2} \\ n_{2} \\ n_{2} \\ n_{3} \\ n_{4} \\ n_$$

METHOD 2.

$$x(n)$$

$$\frac{3 2 (1 2)}{1 (2 + 1)^{2}}$$

$$h(n) = \frac{3}{6} \frac{2}{4} \frac{1}{2} \frac{1}{2}$$

$$h(n) = \frac{3}{6} \frac{2}{4} \frac{1}{2} \frac{1}{4} \frac{2}{4}$$

$$y(n) = \sqrt{3}, 8, 8, 12, 9, 4, 43$$

$$\frac{1}{6} \frac{1}{4} \frac{1}{2} \frac{1}{4}$$

$$y(n) = \sqrt{3}, 8, 8, 12, 9, 4, 43$$

$$\frac{1}{8} \frac{1}{8 - 6}$$
Find the convolution of the signals
$$x(n) = 1 \quad n = -3, 0, 1$$

$$= 2 \quad n = -1$$

$$= 0 \quad \text{elsevature}$$

$$h(n) = \delta(n) - \delta(n-1) + \delta(n-2) + \delta(n-3)$$

$$\text{solution:}$$

$$\frac{1}{3} \frac{1}{-2} \frac{1}{-2} \frac{1}{-1} \frac{1}{6} \frac{1}{8}$$

$$\frac{1}{-3} - 2 - 1 \circ 1 + 2 \cdot 3 \cdot 4 \cdot 5 \cdot 6 \cdot 7$$

$$h(n) = \frac{1}{8} \frac{1}{6} \frac{1}{8} \frac{1}{6} \frac{1}{8} \frac{1}{6} \frac{1}{8} \frac{1}{6} \frac{1}{16} \frac{1}{8} \frac{1}{16} \frac{1}{16}$$

For
$$n_{2-1}$$

 $y_{(-1)} = \sum_{k=\infty}^{n} x_{(k)} h_{(1-k)}$
 $z = 1(-1) + 2(1) = 1$
 $y_{(0)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2+k)}$
 $y_{(0)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2-k)}$
 $y_{(-1)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2-k)}$
 $z = 1(-1) + 2(1) + 1(1) = 0$
For n_{2}
 $y_{(-1)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2-k)}$
 $z = 1(-1) + 2(1) + 1(-1) + 1(1)$
 $z = 1$
for n_{2}
 $y_{(k)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2-k)}$
 $z = 1(-1) + 2(1) + 1(-1) + 1(1)$
 $z = 1$
for n_{2}
 $y_{(k)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2-k)}$
 $z = -2$
For n_{2}
 $y_{(k)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2-k)}$
 $z = -2$
For n_{2}
 $y_{(k)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2-k)}$
 $z = 1(-1) + 1(1)$
 $z = 0$
For n_{2}
 $y_{(k)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2-k)}$
 $y_{(k)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2-k)}$
 $z = 1(-1) + 1(1)$
 $z = 0$
For n_{2}
 $y_{(k)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2-k)}$
 $z = 1(-1) + 1(-1)$
 $z = 0$
For n_{2}
 $y_{(k)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2-k)}$
 $z = 1(-1) + 1(-1)$
 $z = 0$
For n_{2}
 $y_{(k)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2-k)}$
 $z = 1(-1) + 1(-1)$
 $z = 0$
For n_{2}
 $y_{(k)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2-k)}$
 $z = 1(-1) - 1$
 $y_{(k)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2-k)}$
 $z = 1(-1) - 1$
 $y_{(k)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2-k)}$
 $z = 1(-1) - 1$
 $y_{(k)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2-k)}$
 $z = 1(-1) - 1$
 $y_{(k)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2-k)}$
 $z = 1(-1) - 1$
 $y_{(k)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2-k)}$
 $y_{(k)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2-k)}$
 $z = 1(-1) - 1$
 $y_{(k)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2-k)}$
 $z = 1(-1) - 1$
 $y_{(k)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2-k)}$
 $z = 1(-1) - 1$
 $y_{(k)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2-k)}$
 $z = 1(-1) - 1$
 $y_{(k)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2-k)}$
 $z = 1(-1) - 1$
 $y_{(k)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2-k)}$
 $z = 1(-1) - 1$
 $y_{(k)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2-k)}$
 $z = 1(-1) - 1$
 $y_{(k)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2-k)}$
 $z = 1(-1) - 1$
 $y_{(k)} = \sum_{k=\infty}^{n} x_{(k)} h_{(2-k)}$
 $z = 1(-1) - 1$
 $z = 1$
 z

3) Find the convolution of two finite sequences

him = anuin fa all n i'y when a = b

x (n) = bⁿu(n) for all n ily when a 2 b

Solution

The impulse response h 101=0 for nro, so the given bythe is causal and 20120 for nro, hence the sequence is a causal sequence.

 $y(n) = a^n \frac{s}{k_{20}} (1)^k = a^n (n+1)^{-1} + 1 + 1 + 1 + \dots + n + p + errors = n + 1$

4) Find y(n) if z(n)= n+2 for 0≤n≤3 hins an uns for all n Solution :

K=0

$$y(n) = \sum_{k=0}^{\infty} x(k)h(n-k)$$
Given $x(n) = n+2$ for $0 \le n \le 3$
 $h(n) = conu(n)$ for all n
 $f(n) = conu(n)$ for all n
 $f(n) = conu(n)$ for all n
 $x(n)$ is a causal finite sequence value is
zero for $n>8$.
 $y(n) = \sum_{k=0}^{3} x(k)h(n-k)$
 $= \frac{3}{2} (k+3) a^{n-k} y(n-k)$

 $= 2a^{n}u(n) + 3a^{n-1}u(n-1) + 4a^{n-2}u(n-2) + 5a^{n-3}u(n-3)$

5) Determine the response of the relaxed system characterized
by the impute response
$$h(n) = (\frac{1}{2})^n u(n)$$
 to the imput signal
 $X(n) = a^n u(n)$.
Solution:
A causal signal is applied to a causal System.
 $\therefore y(n) = \sum_{k=0}^{n} x(k)h(n-k)$
 $= \sum_{k=0}^{n} x(k)h(n-k)$
 $= \sum_{k=0}^{n} x^k (\frac{1}{2})^{n-k}$
 $= (\frac{1}{4})^n \sum_{k=0}^{n} x^k (\frac{1}{2})^{n-k} = (\frac{1}{2})^n \sum_{k=0}^{n} x^k (\frac{1}{2})^{n-k}$
 $= (\frac{1}{4})^n \sum_{k=0}^{n} x^k (\frac{1}{2})^{n-k}$
 $= (\frac{1}{4})^n \sum_{k=0}^{n} x^k (\frac{1}{2})^{n-k} = (\frac{1}{4})^n \sum_{k=0}^{n-k} x^k (\frac{1}{2})^{n-k}$
 $= (\frac{1}{4})^n \sum_{k=0}^{n-k} x^k (\frac{1}{2})^{n-k} = (\frac{1}{4})^n \sum_{k=0}^{n-k} x^k (\frac{1}{2})^{n-k} = (\frac{1}{4})^n \sum_{k=0}^{n-k} x^k (\frac{1}{2})^{n-k} = (\frac{1}{4})^n \sum_{k=0}^{n-k} x^{k-1} = \frac{1}{4}$
 $= (\frac{1}{4})^n \sum_{k=0}^{n-k} x^{k-1} = \frac{1}{4}$
 $x_1(n) = \int 1, -1, -2, 3, -1 \sum_{k=0}^{n-k} x^{k-1} = \frac{1}{2}$
 $x_1(n) = \int 1, -1, -2, 3, -1 \sum_{k=0}^{n-k} x^{k-1} = \frac{1}{2}$
 $x_1(n) = \int 1, -1, -2, 3, -1 \sum_{k=0}^{n-k} x^{k-1} = \frac{1}{2}$

y (0) = 1(1) + 0 (+) + 0 (2) + 3(3) + 2(+) 2(3) y (0) = 8

$$\begin{split} & y_{1}(1) = 1(2) + (+1)(1) + (+2)(0) + 3(0) + 4(0)(0) \\ & = -4 \\ & y_{1}(3) = 0(1) + 3(-1) + 2(-2) + 1(3) + (+1)(0) \\ & = -4 \\ & y_{1}(3) = 0(1) + 0(-1) + 3(-2) + 3(2) + 3(2) + 1(-1) \\ & = -4 \\ & y_{1}(3) = 0(1) + 0(-1) + 3(-2) + 3(2) + 3(2) + 1(-1) \\ & = -4 \\ & y_{1}(3) = 0(1) + 0(-1) + 3(-2) + 3(2) + 3(2) + 1(-1) \\ & = -4 \\ & y_{1}(3) = 0(1) + 0(-1) + 3(-2) + 3(2) + 3(2) + 1(-1) \\ & = -4 \\ & y_{1}(3) = 0(1) + 0(-1) + 3(-2) + 3(2) + 3(2) + 1(-1) \\ & = -4 \\ & y_{1}(3) = 0(1) + 0(-1) + 3(-2) + 3(2) + 3(2) + 1(-1) \\ & = -1 \\ & y_{1}(3) = 1(1) + 3(0) + 3(2) + 3(0) + 3$$

CORRELATION
Connectation is boostally used to Compare
two Signals
It is a measure of Similarity between
two signals.
Types
is Abtro correlation
It is a measure of Similarity among
the same signal.
The autoconcelation of a Sequence 2000
is defined by

$$N_{HX}(l) = \sum_{n=\infty}^{\infty} 2(m) 2(n-l)$$

 $N_{HX}(l) = \sum_{n=\infty}^{\infty} 1(m) 2(n-l)$
 $N_{HX}(l) = \sum_{n=\infty}^{\infty} 1(m) 2(n-l)$
If the time shift loo, then we have
 $N_{XX}(p) = \sum_{n=\infty}^{\infty} n(p+1) 2(n)$
if cross correlation
It is a measure of Similarity between
two different signals.
The cross concelation between the signal
 $n(n) I y(n)$ is given by
 $N_{XY}(l) = \sum_{n=\infty}^{\infty} 2(n) y(n-l) \rightarrow \otimes l = 0, \pm 1, \pm 2, \dots$
 $N_{YX}(l) = \sum_{n=\infty}^{\infty} y(n+1) 2(n-l)$
 $= \sum_{n=\infty}^{\infty} y(n+1) 2(n-l)$

If the time shift Leo, then we get

$$y_{xy}(o) = y_{yz}(o) = \sum_{n=\infty}^{\infty} x_{(n)} y_{(n)}$$

Comparing equation (a) 2 (b) we find that
 $y_{xy}(t) = y_{yz}(-t)$
where $y_{yz}(-t)$ is the folded version of $y_{yz}(t)$
about Leo.
We can rewrite the equation (b)
 $y_{zy}(t) = \sum_{n=\infty}^{\infty} x_{(n)} y_{(-t)-n} J$
 $= x(t) + y(-t)$
From the above equation, we find that
the correlation process is essentially the involution
of two data sequences in which one of the equanes
has been reversed.
Hind the cross correlation of two first length sequences $x_{(n)} = \int_{1}^{1} x_{y}(t) = \frac{1}{2} \int_{1}^{1} \frac{1}{$

V

UNIT-3

TWIDDLE FACTOR (GR) WINDOW FUNCTION
The twiddle factor or window function in
given rby

$$\begin{split} & W_{N} = e^{-dt} M \\ & W_{N} = e^{-dt$$

FREQUENCY TRANSFORMATION
INTRODUCTION TO DET
The Dissect Fouries Transform (DET) is
a powerful computation tool which follows us to
evaluate the Fouries Transform
$$\chi(e^{j_0})$$
 bu
a digital computer or specially designed hardware.
DET is defined only for sequence of
fruite length.
DET PAIR
The DET of a discrete signal 2000 is given by
 $\chi(e^{j_0})\chi(e^{j_0}) = \sum_{h=0}^{N-1} \chi(n) = \frac{1}{200}$
 $\chi(k) = \chi(e^{j_0})$ bit N equally spaced points
ever $0 \le w \le 2\pi$, we obtain
 $\chi(k) = \chi(e^{j_0}) \int_{0=2\pi k/M}^{0} \sum_{h=0}^{N-1} \sum_{j=2\pi k/M}^{0} \sum_{j=2\pi k/M}^{0}$

Substituting Eqn. (b) is (c) we get

$$x(x) = \sum_{n=0}^{N+1} \left[\frac{1}{N} \sum_{k=0}^{N-1} x(k) e^{\frac{1}{2} 2nk/N} \sum_{k=0}^{n-1} x(k) e^{\frac{1}{2} 2nk/N} \sum_{k=0}^{n-1} x(k) - \frac{1}{2} \sum_{k=0}^{n-1} \frac{x(k)}{1-e^{\frac{1}{2}2nk/N}} \sum_{k=0}^{n-1} \frac{x(k)}{1-e^{\frac{1}{2}2nk/N}} \sum_{k=0}^{n-1} \frac{1}{1-e^{\frac{1}{2}2nk/N}} \sum_{k=0}^{n-1} \frac{1}{1-e^{\frac{1}{2}2n$$

PROOF
DFT
$$[x(N-n)] = \sum_{n=0}^{N-1} x(N-n) e^{-\frac{1}{2}ankn/N}$$

changing the knotex from n to m = N-n we get
 $h = \sum_{n=0}^{N-1} x(m) e^{\frac{1}{2}ank(N-m)/N}$
DFT $[x(N-n)] = \sum_{m=0}^{N-1} x(m) e^{\frac{1}{2}ank(N-m)/N}$
 $= \sum_{m=0}^{N-1} x(m) e^{\frac{1}{2}anm(N-n)/N}$
 $N = \sum_{m=0}^{N-1} x(m) e^{\frac{1}{2}anm(N-n)/N}$
 $N = \sum_{m=0}^{N-1} x(m) e^{\frac{1}{2}anm(N-n)/N}$
 $PT [x(n) e^{\frac{1}{2}an(n/n)}] = x((k-1))N$
PROOF!
DFT $[x(n) e^{\frac{1}{2}an(n/n)}] = x((k-1))N$
PROOF!
DFT $[x(n) e^{\frac{1}{2}an(n/n)}] = \sum_{m=0}^{N-1} x(m) e^{\frac{1}{2}ann(N-n)/N} e^{\frac{1}{2}ann/N}$
 $x(n) = [x(n), x(n), ..., x(n-1)] = \sum_{m=0}^{N-1} x(m) e^{\frac{1}{2}ann(N-n)/N}$
 $x(n-n)_{N} = [x(n, n, ..., x(n-1)] = \sum_{m=0}^{N-1} x(m) e^{-\frac{1}{2}ann(N-n-1)/N}$
 $x(n-n)_{N} = [x(n, n, ..., x(n-1)] = \sum_{m=0}^{N-1} x(m) e^{-\frac{1}{2}ann(N+n-1)/N}$
 $x(n-n)_{N} = [x(n-n)_{N} = x(n-n-1)] = x((n-n)/N$
 $x(n-n)_{N} = x(n-n-1) = x((n-1)/N$
 $N = \sum_{m=0}^{N-1} x(m) = x((n-1)/N$
 $N = \sum_{m=0}^{N-1} x(m) = \sum_{m=0}^{N-1} x(m) e^{-\frac{1}{2}ann(N+n-1)/N}$
 $N = \sum_{m=0}^{N-1} x(m-n-1) = x((n-n-1)/N$
 $N = \sum_{m=0}^{N-1} x(m-n-1) = x((n-n-1)/N$
 $N = \sum_{m=0}^{N-1} x(m-n-1) = x((n-n-1)/N$
 $N = \sum_{m=0}^{N-1} x(m-1) = x((n-1)/N$
 $N = \sum_{m=0}^{N-1} x(m-1) = x((n-1)/N$
 $N = \sum_{m=0}^{N-1} x(m-1) = x((n-1)/N$
 $N = \sum_{m=0}^{N-1} x(m-1) = x((n-$

PROOF
DFT [
$$x^{*}(n)^{T}$$
] = $\sum_{n=0}^{M-1} x^{*}(n) = \int_{a \in M^{n}/M}^{a \times M^{n}/M}$
= $\left[\sum_{n=0}^{M-1} x(n) = \int_{a \in M^{n}/M}^{a \times M^{n}/M}\right]^{*}$
= $\left[\sum_{n=0}^{M-1} x(n) = \int_{a \in M^{n}/M}^{a \times M^{n}/M}\right]^{*}$
DFT [$x^{*}(n) = \int_{a \in M^{n}}^{M-1} x^{*}(n) = \int_{a \in M^{n}/M}^{a \times M^{n}/M}$
DFT [$x^{*}(n) = \int_{a \in M^{n}/M}^{M-1} x^{*}(n) = \int_{a \in M^{n}/M}^{a \times M^{n}/M}$
= $\int_{a \in M^{n}/M}^{M-1} x(n) = \int_{a \in M^{n}/M}^{a \times M^{n}/M}$
= $\int_{a \in M^{n}/M}^{M-1} x(n) = \int_{a \in M^{n}/M}^{M-1} x(n) = \int_{a \in M^{n}/M}^{a \times M^{n}/M}$
= $\int_{a \in M^{n}/M}^{M-1} x(n) = \int_{a \in M^{n}/M}^{M-1} x(n) =$

DFT
$$[x_{i}(n) \oplus x_{2}(n)] = x_{i}(k) x_{2}(k)$$

Viis circular correlation
For complex valued sequences $x(n)$ and $y(n) + i$
DFT $[x(n)] = x(k)$ and
DFT $[y(n)] = y(k)$.
then
DFT $[\overline{y}_{ny}(l)] = DFT \begin{bmatrix} N^{-1} - x(n) y^{*}((n-1)) \\ n = a \\ m = a \\$

0 Determine the 4-Detrit DFT and IDFT of the given signal

$$u(n) = \begin{cases} 1, & 0 \le n \le 3 \\ 0, & \text{elemetric} \end{cases}$$

Sol:
The ilp Sequence is $u(n) = \int [1, 1, 1, 1]$. Do this problem
 $L = N = 4$.
DFT of $u(n)$ is, $x(n) = \sum_{\substack{n=0 \\ n \ge 0}}^{N-1} \sum_{\substack{n=0 \\ n \ge 0}}^{-j \ge n} nh$
 $u(n) = \sum_{\substack{n=0 \\ n \ge 0}}^{3} u(n) e^{-j \underbrace{n} + nh}$
 $u(n) = \sum_{\substack{n=0 \\ n \ge 0}}^{j} u(n) e^{-j \underbrace{n} + nh}$
 $u(n) = \sum_{\substack{n=0 \\ n \ge 0}}^{j} u(n) e^{-j \underbrace{n} + nh}$
 $u(n) = \sum_{\substack{n=0 \\ n \ge 0}}^{j} u(n) e^{-j \underbrace{n} + nh}$
 $u(n) = \sum_{\substack{n=0 \\ n \ge 0}}^{j} u(n) e^{-j \underbrace{n} + nh}$
 $u(n) = \sum_{\substack{n=0 \\ n \ge 0}}^{j} u(n) e^{-j \underbrace{n} + nh}$
 $u(n) = \frac{1}{2} u(n) e^{-j \underbrace{n} + nh}$
 $u(n) = 1 + 1 + 1 + 1 = 4$
For $k \ge 1$
 $u(n) = \sum_{\substack{n=0 \\ n \ge 0}}^{j} \underbrace{n \le 0}_{n \ge 0} e^{-j \underbrace{n} + nh}_{n \ge 0} e^{-j \underbrace{n} + u(n)}_{n \ge 0} e^{-j \underbrace{n} \underbrace{n} e^{-j \underbrace{n} + u(n)}_{n \ge 0} e^{-j \underbrace{n} + u(n)}_{n \ge 0} e^{-j \underbrace{n} \underbrace{n} \underbrace{n} e^{-j \underbrace{n} + u(n)}_{n \ge 0} e^{-j \underbrace{n} \underbrace{n} \underbrace{n} \underbrace{n} e^{-j \underbrace{n} e^{-j \underbrace{n} + u(n)}_{n \ge 0} e^{-j \underbrace{n} \underbrace{n} e^{-j \underbrace{n} \underbrace{n} \underbrace{n} e^{-j \underbrace{n} \underbrace{n} \underbrace{n} \underbrace{n} e^{-j \underbrace{n} \underbrace{n} \underbrace{n} \underbrace{n} \underbrace$

$$= 1 + (e^{s_{1}} - j^{s_{1}} e^{s_{1}} - j^{s_{1}} e^{s_{1}} + (e^{s_{2}} - j^{s_{1}} e^{s_{1}}) + (e^{s_{2}} - j^{s_{1}} e^{s_{1}})$$

$$= 1 + (e^{-j}) + (-1 - j^{s_{1}}) + (e^{-j})^{s_{1}}$$

$$= 1 + (e^{-j}) + (1 - e^{-j}) + (e^{-j})^{s_{1}}$$

$$= 1 + (e^{-j}) + (1 - e^{-j}) + (e^{-j})^{s_{1}}$$

$$x(e) = \frac{3}{2} + 2(e^{-j})^{s_{1}} e^{s_{1}} + 2(e^{-j})^{s_{1}} + 2(e^{-j})^{s_{1}} e^{s_{1}}$$

$$= 1 + (e^{s_{1}})^{s_{1}} e^{-j^{s_{1}}} e^{s_{1}} + 2(e^{j})e^{-j^{s_{1}}} + 2(e^{-j})^{s_{1}} e^{-j^{s_{1}}} e^{s_{1}}$$

$$= 1 + (e^{s_{1}})^{s_{1}} e^{-j^{s_{1}}} e^{s_{1}} e^{s_{1}} e^{s_{1}} e^{s_{1}} e^{-j^{s_{1}}} e^{-j^{s_{1}}}} e^{-j^{s_{1}}} e^{-j^{s_{1}}} e^{-j^{s_{1}}}} e^{-j^{s_{1}}} e^{-j^{s_{1}}} e^{-j^{s_{1}}} e^{-j^{s_{1}}} e^{-j^{s_{1}}} e^{-j^{s_{1}}} e^{-j^{s_{1}}}} e^{-j^{s_{1}}} e^{-j^{s$$
For
$$n \ge 1$$

 $\pi(t) = \frac{1}{4} \sum_{k=0}^{3} \chi(k) e^{j \frac{\pi}{2}k}$, $n \ge 0, 1, 2, 3$
 $= \frac{1}{4} \left[\chi(0) e^{0} + \chi(t) e^{j \frac{\pi}{2}} + \chi(2) e^{j \pi} + \chi(3) e^{j \frac{\pi}{2}} \right]$
 $= \frac{1}{4} \left[4 + 0 + 0 + 0 \right]$

2602 = 1

For
$$n = 2$$
,
 $2(2) = \frac{1}{4} \sum_{k=0}^{3} x(k) e^{j\pi k}$, $n = 0, 1, 2, 3$
 $2(2) = \frac{1}{4} [x 0) e^{0} + x(0) e^{j\pi} + x(2) e^{j2\pi} + x(3) e^{j3\pi}]$
 $= \frac{1}{4} [4 + 0 + 0 + 0]$

21(2) = 1

For
$$n=3$$

 $n(3) = \frac{1}{4} \begin{bmatrix} 3 \\ k > 0 \end{bmatrix} \begin{bmatrix} 3 \\ k > 0 \end{bmatrix} \begin{bmatrix} 3 \\ k > 0 \end{bmatrix} \begin{bmatrix} 1 \\ k > 0 \end{bmatrix} \begin{bmatrix} 3 \\ k > 0 \end{bmatrix} \begin{bmatrix} 3 \\ k > 0 \end{bmatrix} \begin{bmatrix} 1 \\ k > 0 \end{bmatrix} \begin{bmatrix} 2 \\ k > 0 \end{bmatrix} \begin{bmatrix} 3 \\ k > 0 \end{bmatrix} \begin{bmatrix}$

2(3)=1 i. The IDET is , 2007={1,1,1,1}

Find the DFT of a Sequence
$$x(0) = \{1, 1, 0, 0\}$$
 and find
the IDFT of $\gamma(K) = \{1, 0, 1, 0\}$
Solution
Let us assume $N = L = A$
DFT of $x(0)$ is given as
 $x(K) = \sum_{n=0}^{M} x(0) e^{-j\alpha nK/N}$. $K = 0, 1, ..., N = 1$.
Kso:
 $x(0) = \sum_{n=0}^{M} x(0) = x(0) + x(0) + x(0) + x(0) + x(0)$
 $= 1 + 1 + 0 + 0$
 $= 2$
 $K = 1$
 $x(1) = \sum_{n=0}^{M} x(0) e^{-jn/2} = x(0) + x(0) e^{-jn/2} + x(0) e^{-jn/2} + x(0) e^{-jn/2}$
 $= 1 + \cos \frac{\pi}{2} - \int \sin \frac{\pi}{2}$
 $K = R$
 $K = R$
 $K = 2$
 $K = 3$
 $x(2) = \sum_{n=0}^{M} x(0) e^{-jn/2} = x(0) + x(0) e^{-jn} + x(0) e^{-jn} + x(0) e^{-jn}$
 $= 1 + \cos \frac{\pi}{2} - \int \sin \frac{\pi}{2}$
 $K = 3$
 $= 1 + \cos \frac{\pi}{2} - \int \sin \frac{\pi}{2}$
 $K = 3$
 $= 1 + \cos \frac{\pi}{2} - \int \sin \frac{\pi}{2} + x(0) e^{-jn} + x(0) e^{-jn} + x(0) e^{-jn}$
 $= 1 + \cos \frac{\pi}{2} - \int \sin \frac{\pi}{2} + x(0) e^{-jn} + x(0) e^{-jn} + x(0) e^{-jn} + x(0) e^{-jn}$
 $= 1 + \cos \frac{\pi}{2} - \int \sin \frac{\pi}{2} + x(0) e^{-jn} + x(0) e^{$

T.

$$\begin{array}{l} y(1) = \frac{1}{N} \frac{3}{K_{0}} y(k) e^{j\pi k/2} \\ = \frac{1}{N} \frac{1}{K_{0}} y(k) e^{j\pi k/2} \\ = \frac{1}{N} \frac{1}{K} (y(0) + y(1)) e^{j\pi k} + y(2) e^{j\pi k} + y(3) e^{j\pi k} \end{bmatrix} \\ = \frac{1}{N} \frac{1}{1+0} + \cos \pi + j\sin \pi + 0 \end{bmatrix} \\ = \frac{1}{N} \frac{1}{1+0} + 1+0 \end{bmatrix} = 0 . \\ y(2) = \frac{1}{4} \frac{1}{N} (y(0) + y(1)) e^{j\pi k} + y(3) e^{j\pi \pi} + y(3) e^{j\pi \pi} \end{bmatrix} \\ = \frac{1}{N} \frac{1}{1+0} + 1+0 \end{bmatrix} = 0 . \\ y(3) = \frac{1}{N} \frac{1}{N} (y(0) + y(1)) e^{j\pi k/2} + y(3) e^{j\pi k/2} \end{bmatrix} \\ = \frac{1}{N} \frac{1}{N} \frac{1}{N} + y(3) e^{j\pi k/2} + y(3) e^{j\pi k/2} \end{bmatrix} \\ = \frac{1}{N} \frac{1}{N} \frac{1}{N} + y(3) e^{j\pi k/2} + y(3) e^{j\pi k/2} \end{bmatrix} \\ = \frac{1}{N} \frac{1}{N$$

For N=4

$$x(k) = \sum_{n=0}^{3} x(n) e^{-jx_{n}k/2}$$

 $k=0, 1, 2, 3$
 $k=0$
 $x(0) = x_{0}(0) + x(0) e^{jx_{2}} + x(2) + x(3)$
 $= 3$
 $k=1$
 $x(1) = x(0) e^{-0} + x(1) e^{jx_{2}} + x(2) e^{-jx_{n}} + x(3) e^{-jx_{n}}$
 $= 1 + \cos \frac{\pi}{2} - j\sin \frac{\pi}{2} + x(2) e^{-\frac{\pi}{2}} + x(3) e^{-\frac{\pi}{2}} + x(3) e^{-\frac{\pi}{2}}$
 $= 1 + \cos \frac{\pi}{2} - j\sin \frac{\pi}{2} + x(2) e^{-\frac{\pi}{2}} + x(3) e^{-\frac{\pi}{2}} + x(3) e^{-\frac{\pi}{2}}$
 $k = 3$
 $x(3) = x(0) + x(1) e^{jx_{1}} + x(2) e^{-\frac{\pi}{2}} + x(3) e^{-\frac{\pi}{2}} + x(3) e^{-\frac{\pi}{2}}$
 $= 1 + \cos \frac{\pi}{2} - j\sin \frac{\pi}{2} + \cos \frac{\pi}{2} - j\sin \frac{\pi}{2} + \cos \frac{\pi}{2}$

For
$$k=1$$

 $x(t) = \int_{1-\infty}^{T} x(t)e^{-jxt}/4$
 $= \pi(0) + \pi(t)e^{-j\pi/4} + \pi(2)e^{-j\pi/2}$.
 $= 1+0.707 - j0.707 + 0-j$
 $= 1.707 - j1.707$
For $k=2$
 $\pi=0$
 $= 1.707 - j1.707$
For $k=3$
 $x(2) = \int_{1-\infty}^{\infty} \pi(t)e^{-j\pi/4} + \pi(2)e^{-j\pi}$
 $= 1 + \cos \frac{\pi}{2} - j\sin \frac{\pi}{4} + \pi(2)e^{-j\frac{\pi}{2}\pi/2}$
 $= 1 + \cos \frac{\pi}{2} - j\sin \frac{\pi}{4} + \pi(2)e^{-j\frac{\pi}{2}\pi/2}$
 $= 1 + \cos \frac{\pi}{4} - j\sin \frac{\pi}{4} + \pi(2)e^{-j\frac{\pi}{4}\pi/2}$
 $= \pi(e) + \pi(t)e^{-j\frac{\pi}{4}} + \pi(2)e^{-j\frac{\pi}{4}\pi/2}$
 $= 1 + \cos \frac{\pi}{4} - j\sin \frac{\pi}{4} + \pi(2)e^{-j\frac{\pi}{4}\pi/2}$
 $= 1 + \cos \frac{\pi}{4} - j\sin \frac{\pi}{4} + \cos \frac{\pi}{4} - j\sin \frac{\pi}{4}$
 $= \pi(e) + \pi(t)e^{-j\frac{\pi}{4}} + \pi(2)e^{-j\frac{\pi}{4}\pi/2}$
 $= 1 + \cos \frac{\pi}{4} - j\sin \frac{\pi}{4} + \cos \frac{\pi}{4} - j\sin \frac{\pi}{4}$
 $= \pi(0) + \pi(t)e^{-j\frac{\pi}{4}} + \pi(2)e^{-j\frac{\pi}{4}\pi/2}$
 $= \pi(0) + \pi(t)e^{-j\frac{\pi}{4}} + \pi(2)e^{-j\frac{\pi}{4}\pi/2}$
 $= \pi(0) + \pi(t)e^{-j\frac{\pi}{4}} + \pi(2)e^{-j\frac{\pi}{4}\pi/2}$
 $= \pi(e) + 2(t)e^{-j\frac{\pi}{4}} + \pi(2)e^{-j\frac{\pi}{4}\pi/2}$
 $= \pi(e) + 2(t)e^{-j\frac{\pi}{4}} + \pi(2)e^{-j\frac{\pi}{4}\pi/2}$
 $= \pi(e) + \pi(1)e^{-j\frac{\pi}{4}} + \pi(2)e^{-j\frac{\pi}{4}\pi/2}$
 $= 1 + \cos \pi - j\sin \pi + \cos \pi - j\sin \pi + \cos \pi - j\sin \pi + \cos \pi/2 + \sin \pi/2$
 $= \pi(e) + \pi(1)e^{-j\frac{\pi}{4}} + \pi(2)e^{-j\frac{\pi}{4}\pi/2}$
 $= \pi(e) + \pi(1)e^{-j\frac{\pi}{4}} + \pi(2)e^{-j\frac{\pi}{4}\pi/2}$
 $= 1 + \cos \frac{\pi}{4} - j\sin \frac{\pi}{4} + \cos \frac{\pi}{4} + \cos \frac{\pi}{4} - \sin \frac{\pi}{4}$
 $= 1 + \cos \frac{\pi}{4} - \sin \frac{\pi}{4} + \cos \frac{\pi}{4} + \cos \frac{\pi}{4} - \sin \frac{\pi}{4}$
 $= 1 + \cos \frac{\pi}{4} - \sin \frac{\pi}{4} + \cos \frac{\pi}{4} + \cos \frac{\pi}{4} - \sin \frac{\pi}{4}$

For k=6

$$x(b) = \int_{a=0}^{b} x(0) e^{-jB\pi n/2}$$

 $= 2x(0) + 2x(0) e^{-jB\pi n/2}$
 $= 1x(0) = \int_{a=0}^{b} x(0) e^{-jB\pi n/2}$
 $= 1x(0) e^{-jB\pi n/2}$
 $= 1x(0) = \int_{a=0}^{b} x(0) e^{-jB\pi n/2}$
 $= 1x(0) e^{-jB\pi n/2}$
 $x(0) e^{-jB\pi n/2}$
 $x(0) e^{-jB\pi n/2}$
 $= 1x(0) e^{-jB\pi n/2}$
 $x(0) e^{-jB\pi n/2}$
 $= 1x(0) e^{-jB\pi n/2}$
 $x(0) e^{-jB\pi n/2}$
 $= 1x(0) e^{-jB\pi n/2}$
 $x(0) e^{-jB\pi n/2}$
 $x(0) e^{-jB\pi n/2}$
 $= 1x(0) e^{-jB\pi n/2}$
 $x(0) e^{-jB\pi n/2}$
 $= 1x(0) e^{-jB$

For n=1

$$x(1) = \frac{1}{8} \left[\frac{7}{k=0} \times (k) e^{j\pi k/4} \right]$$

 $= \frac{1}{8} \left[5 + (1-j) j + 1(-1) + (1+j) (j) \right]$
 $= \frac{1}{8} \left[5 + (1-j) (j) + 1(-1) + (1+j) (-1) \right]$
 $= \frac{1}{8} \left[5 + (1-j) (-1) + 1(-1) + (1+j) (-1) \right]$
 $= \frac{1}{8} \left[4 \right] = 0.5$
For n=3
 $7(8) = \frac{1}{8} \left[\frac{7}{k=0} \times (k) e^{j\pi k/4} \right]$
 $= \frac{1}{8} \left[5 + (1-j) (j) + 1(-1) + (1+j) (j) \right]$
 $= \frac{1}{8} \left[2 - 0.3 e^{-5} \right]$
For n=4
 $\pi(4) = \frac{1}{8} \left[\frac{7}{k=0} \times (k) e^{j\pi k/4} \right]$
 $= \frac{1}{8} \left[5 + (1-j) (1) + 1(1) + (1+j) (1) \right]$
 $= 1$
For n=5
 $\pi(5) = \frac{1}{8} \left[\frac{7}{k=0} \times (k) e^{j5\pi k/4} \right]$
 $= \frac{1}{8} (5) = 0.75$
For n=6
 $\pi(6) = \frac{1}{8} \left[\frac{7}{k=0} \times (k) e^{j3\pi k/4} \right]$
 $= \frac{1}{8} \left[5 + (1-j) (2) + 1(1) + (1+j) (-1) \right]$
 $= \frac{1}{8} \left[5 + (1-j) (2) + 1(1) + (1+j) (-1) \right]$
 $= \frac{1}{8} \left[5 + (1-j) (2) + 1(1) + (1+j) (-1) \right]$
For n=6
 $\pi(6) = \frac{1}{8} \left[\frac{7}{k=0} \times (k) e^{j3\pi k/4} \right]$
 $= \frac{1}{8} \left[5 + (1-j) (-1) + 1(1) + (1+j) (-1) \right]$
 $= \frac{1}{8} \left[5 + (1-j) (-1) + 1(1) + (1+j) (-1) \right]$
 $= \frac{1}{8} \left[5 + (1-j) (-1) + 1(1) + (1+j) (-1) \right]$
 $= \frac{1}{8} \left[5 + (1-j) (-1) + 1(1) + (1+j) (-1) \right]$
 $= \frac{1}{8} \left[5 + (1-j) (-1) + 1(1) + (1+j) (-1) \right]$
 $= \frac{1}{8} \left[5 + (1-j) (-1) + 1(1) + (1+j) (-1) \right]$
 $= \frac{1}{8} \left[\frac{7}{k=0} \times (k) e^{j\pi k/4} \right] = 0.85$
For n=7
For n=7
 $\pi(7) = \frac{1}{8} \left[\frac{7}{k=0} \times (k) e^{j\pi k/4} \right] = 0.85$

CIRCULAR. CONVOLUTION BASED ON DET AND IDET METHOD " Perform the cucular convolution of the following sequences 2(n)= {1,1,2,1} & h(n)= {1,2,3,4} using OFT and IDFT nethod Solution. By equation of Cercular Convolution, we have $\gamma(\kappa) = \chi(k) H(\kappa).$ $\chi(k) = \sum_{n=0}^{N-1} \chi(n) e^{-j 2\pi n k}$, k = 0, 1, ... N - 1. Here N=4. 2 x(0) = {1,1,2,1} K=0,1,2,3 For K20. $X(0) = \frac{3}{2} X(0) e^{0}$ = x(0) + x(1) + x(2) + x(3) = 5 . . . Store and Store and a const-Far K=1 $X(1) = \frac{3}{2} \alpha(n) e^{-j2\pi n/4}$ $= x(0) + x(1) e^{-\chi_0} + x(2) e^{-\chi} + x(3) e^{-\frac{1}{3}\chi_0}$ = 1-j-2+j=-1. For 1522 $x(z) = \frac{3}{2} x(n) e^{-j a \pi n}$ = 3 x (0) + x (1) e + x (2) e + x (3) e - j3 x. For k=3 $x(3) = \frac{3}{5} \pi(n) e^{-j3\pi n/3}$ = 21(0) + 21(1) e + 21(2) e + 21(3) e - 19 1/2. -1+J-2-J=-1.

By IDFT we have

$$y_{1}(n) = \frac{1}{N} \sum_{k=0}^{N-1} y(k) e^{j2\pi n k} / n \cdot \frac{1}{N} = 0, 1, \dots N-1,$$

For $h = 0 \cdot \frac{3}{4} = y(k) e^{j2\pi n k} / n \cdot \frac{1}{4} = \frac{1}{4} \int y(0) + y(0) + y(0) + y(0) \int y(0) = \frac{1}{4} = \frac{3}{4} \int y(0) + y(0) + y(0) + y(0) \int y(0) = \frac{1}{4} \int \frac{3}{4} \int y(0) + y(0) + y(0) + y(0) \int \frac{1}{4} \int \frac{3}{4} \int y(0) + y(0) e^{j\pi k} + y(0$

FILTERING METHODS BASED ON DET When the input sequence is of long duration has to be processed is a system whose imposhe and response is of finite duration, then it would not be able to store the data before convolution. Therefore, the input sequence nust be divided I and much the descende into blocks. LAX PAL XX MALANT Each individual blocks are processed Separately one at a time and the results are combined to yeard the destried output sequence which is identical to the sequence obtained by linear convolution. Two methods that are commonly used for fithering the sectioning data and combining the results are the overlap. Save nethod and the overlap add nethod. * Length of din ip sequence be is. + Length of an impube response is M. * Input sequence is devided into blocks of data "Singe 12, 11, 10, 2, 0, 0, 0, 2, ..., 00, 2, 00, 2 = 0.0 12 N=L+H-1. * Each block consists of Last (H-1) data points of Previous block followed by L new data points to form a data sequence of tength N=++N=1. TFor first block of data the first M-1 points are set to zoro.

SECTIONED CONVOLUTION [Filturing Long duration Sequences]. 1) Determine the output of the linear FIR filter whose inepulse response him = f1, 2, 33 and the input signal xin= f1, 2, 3, 4, 5, 6, 7, 8, 9} using overlap add method." Solution ! Given ac(n)= \$ 1,2,3,4,5,6,7,8,9 } < 15 (length of ip sequence) h (n) = {1, 2, 3} = 4 (Length of impuls) sequence y (n) + ++++++ Subdivide the ofp data sequence as 3 because the length of the impube response h(n). To create, a new length of each sequence, N = L + M - 1e et et s fermit = 3+3-1 N=5 To make the length as 5, add (M-1) xerves at each end of the sequence. $x_{1}(0) = \{1, 2, 3, 0, 0\}$ x2(n)= {4,5,6,00 } 23(17)= 57,8,9,0,03. Now, we have if p data sequence with tength 5. So we must make the length of h(n) on 5 by appending M-1 Xeroes. 50 Rm= f1, 2, 3,0,0 }. Y (m) = 201(m) + h m) $\begin{bmatrix} 1 & 0 & 0 & 3 & 9 \\ 2 & 1 & 0 & 0 & 3 \\ 3 & 2 & 1 & 0 & 0 \\ 0 & 3 & 2 & 1 & 0 \\ 0 & 3 & 2 & 1 & 0 \end{bmatrix} \begin{bmatrix} 1 \\ 2 \\ 3 \\ 3 \end{bmatrix} = \begin{bmatrix} 4 \\ 10 \\ 10 \end{bmatrix}$

$$Y_{3}(n) = x_{3}(n) * h(n)$$

$$= \begin{bmatrix} 4 & 0 & 0 & 5 & 5 \\ 0 & 5 & 4 & 0 & 0 \\ 0 & 0 & 5 & 5 & 4 \\ 0 & 0 & 5 & 5 & 4 \\ 0 & 0 & 5 & 5 & 4 \\ 0 & 0 & 5 & 5 & 4 \\ 0 & 0 & 5 & 5 & 4 \\ 0 & 0 & 5 & 5 & 4 \\ 0 & 0 & 5 & 5 & 4 \\ 0 & 0 & 5 & 5 & 4 \\ 0 & 0 & 0 & 9 & 8 \\ 0 & 0 & 9 & 8 & 7 & 0 \\ 0 & 0 & 0 & 9 & 8 \\ Y_{3}(n) = 1 & 4 & 10 & 12 & 9 \\ Y_{3}(n) = 1 & 4 & 10 & 12 & 9 \\ Y_{3}(n) = 1 & 4 & 10 & 12 & 9 \\ Y_{3}(n) = 1 & 4 & 10 & 12 & 9 \\ Y_{3}(n) = 1 & 1 & 0 & 12 & 9 \\ Y_{3}(n) = 1 & 4 & 10 & 12 & 9 \\ Y_{3}(n) = 1 & 1 & 0 & 12 & 9 \\ Y_{3}(n) = 1 & 1 & 0 & 12 & 9 \\ Y_{3}(n) = 1 & 1 & 0 & 12 & 9 \\ Y_{3}(n) = 1 & 1 & 0 & 12 & 9 \\ Y_{3}(n) = 1 & 1 & 0 & 12 & 9 \\ Y_{3}(n) = 1 & 1 & 0 & 12 & 9 \\ Y_{3}(n) = 1 & 1 & 0 & 12 & 9 \\ Y_{3}(n) = 1 & 1 & 0 & 12 & 9 \\ Y_{3}(n) = 1 & 1 & 0 & 12 & 9 \\ Y_{3}(n) = 1 & 1 & 0 & 12 & 9 \\ Y_{3}(n) = 1 & 1 & 0 & 12 & 9 \\ Y_{3}(n) = 1 & 1 & 0 & 12 & 9 \\ Y_{3}(n) = 1 & 1 & 0 & 12 & 9 \\ Y_{3}(n) = 1 & 1 & 0 & 12 & 10 \\ Y_{3}(n) = 1 & 1 & 0 & 12 & 10 \\ Y_{3}(n) = 1 & 1 & 0 & 12 & 10 \\ Y_{3}(n) = 1 & 1 & 0 & 12 & 10 \\ Y_{3}(n) = 1 & 1 & 0 & 12 & 10 \\ Y_{3}(n) = 1 & 1 & 0 & 12 & 10 \\ Y_{3}(n) = 1 & 1 & 0 & 12 & 10 \\ Y_{3}(n) = 1 & 1 & 0 & 12 & 10 \\ Y_{3}(n) = 1 & 1 & 0 & 12 & 10 \\ Y_{3}(n) = 1 & 1 & 0 & 12 & 10 \\ Y_{3}(n) = 1 & 1 & 0 & 12 & 10 \\ Y_{3}(n) = 1 & 1 & 0 & 12 & 10 \\ Y_{3}(n) = 1 & 1 & 0 & 12 & 10 \\ Y_{3}(n) = 1 & 1 & 0 & 12 \\ Y_{3}(n) = 1 & 1 & 0 & 12 \\ Y_{3}(n) = 1 & 1 & 0 & 12 \\ Y_{3}(n) = 1 & 1 & 0 & 12 \\ Y_{3}(n) = 1 & 1 & 0 & 12 \\ Y_{3}(n) = 1 & 1 & 0 & 10 \\ Y_{3}(n)$$

.

Y, (n) =
$$\{-1, 0, 3, 2, 2\}$$
.
Yaln) = $\{-1, 0, 4, 6\}$
Yaln) = $\{-1, 0, 1, 3, 4\}$
.
Hum
Find the off yins of a fitter where impube response is
h on = $\{-1, 1, 1\}$ a typ signal $x(n) = \{-3, -1, 0, 1, 3, 2, 0, 1, 2, -1\}$ write
Overlap some & overlap add method.
H.N:
For the $2t(n), 3t_{2}(n)$ and N given Compute $x(n)$ (D) of n
 $x_{1}(n) = 5(n) + 5(n + 1) + 5(n - 2)$,
 $x_{2}(n) = 25(n) - 5(n - 1) + 25(n - 2)$ $N = 4$. $N = 3$.

FAST FOURIER TRANSFORM ALGORITHM (FFT)

The Fast Fourier Transform (FFT) is a highly efficient procedure for computing the DFT of a finite Services and requires less number of computations than that of direct evaluation of DFT.

FFT computation techniques is used in degital spectral analysis, filter simulation, auto correlation and pattern recognition.

FFT impeoves the performance by a factor los or more over direct evaluation of the DFT.

No. of complex nultiplication in DFT is N2 No, of complex Multiplication in FFT is N log_N.

DECIMATION IN TIME ALGORITHM

This algorithm is also known as Radix-2 DIT FFT algorithm.

The No. of output points N=2", where M is an integer

an integer. Let reins is an N-point sequence, decrimate this sequence into two sequence of length NS, where One sequence consusting of the even-indexed values of non and the other of odd indexed values of zins.

 $i e > ne(n) = \pi(an) \quad n = 0, 1, \dots, \frac{N}{2} - 1$ $x_0(n) = x(en+1)$ $n = 0, 1, n = 1, \frac{N}{2} - 1$

The N-point DFT of alos can be written as $x(h) = \sum_{n=0}^{N+1} x(n) W_N^{nk} K = 0, 1, \dots N-1$

Separating
$$\chi(n)$$
 into even and odd indexed values of
 $\chi(n)$, we obtain
 $\chi(k) = \sum_{D=0}^{N-1} \chi(n) W_N^{nk} + \sum_{D=0}^{N-1} \chi(n) W_N^{nk}$
(even) (even)
(even) (even)
 $\chi(k) = \sum_{D=0}^{N-1} \chi(n) W_N^{nk} + \sum_{D=0}^{N-1} \chi(n+1) W_N^{n+1} + \sum_{D=0}^{N-1} \chi(n+1) W_N^{n+1}$
 $= \sum_{D=0}^{N-1} \chi(n) W_N^{nk} + \sum_{D=0}^{N-1} \chi(n+1) W_N^{nk}$
 $= \sum_{D=0}^{N-1} \chi(n) W_N^{nk} + \sum_{D=0}^{N-1} \chi(n+1) W_N^{nk}$
 $= \sum_{D=0}^{N-1} \chi(n) W_N^{nk} + W_N^{nk} \sum_{D=0}^{N-1} \chi(n+1) W_N^{nk}$
The above equation is in terms of even and odd
indexed values of $\chi(n)$:
 $\chi(h) = \sum_{D=0}^{N-1} \chi_0(n) W_N^{nk} + W_N^{nk} \sum_{D=0}^{N-1} \chi_0(n) W_N^{nk}$
 $W_N^{nk} = \left[e^{-i\pi \chi_N}\right]^2 = e^{-i\pi \chi_N/2}$
 $W_N^{nk} = \left[e^{-i\pi \chi_N}\right]^2 = e^{-i\pi \chi_N/2}$
 $W_N^{nk} = \sum_{D=0}^{N-1} \chi_0(n) W_N^{nk} + W_N^{nk} \sum_{D=0}^{N-1} \chi_0(n) W_N^{nk}$
 $\chi(h) = \sum_{D=0}^{N-1} \chi_0(n) W_N^{nk} + W_N^{nk} \sum_{D=0}^{N-1} \chi_0(n) W_N^{nk}$
 $\chi(h) = \sum_{D=0}^{N-1} \chi_0(n) W_N^{nk} + W_N^{nk} \sum_{D=0}^{N-1} \chi_0(n) W_N^{nk}$
 $\chi(h) = \chi_0(h) + W_N^{nk} \times \chi_0(h)$ $H=0, 1, \dots, M-1$
 $\chi_0(h) = \chi_0(h) + W_N^{nk} \times \chi_0(h)$ $H=0, 1, \dots, M-1$



DECIMATION IN FREQUENCY ALGORITHM

In DIF algorithm the output sequence x(k) is divided into smaller and smaller subsequences.

In this algorithm the input sequence ain) is partitioned into two sequences each of length Mg Samples.

The first sequence xi(n) consist of first N/2 Samples of xin) and the sequence x2(n) consists of the last N/2 samples of xin).

$$p_{e_1}(n) = q_{e_1}(n)$$
, $n = 0, 1, d, \dots, N_d - 1$

$$\chi_{a}(n) = \chi(n+N_{a}), n=0,1,a,\dots,N_{a}-1$$

If N=8, x,(n) has values for 0≤n≤3 and 2(n) has values for 4≤n≤7.

The N-point DFT of
$$x(n)$$
 can be wouthen as
 $x(k) = \int_{n=0}^{N-1} x(n) W_N^{nk} + \int_{n=0}^{N-1} x(n) W_N^{nk}$
 $= \int_{n=0}^{N-1} x(n) W_N^{nk} + \int_{n=0}^{N-1} x_2(n) W_N^{nk}$
 $= \int_{n=0}^{N-1} x_1(n) W_N^{nk} + W_N^{nk} \int_{n=0}^{N-1} x_2(n) W_N^{nk}$
 $= \int_{n=0}^{N-1} x_1(n) W_N^{nk} + W_N^{nk} \int_{n=0}^{N-1} x_2(n) W_N^{nk}$
 $= \int_{n=0}^{N-1} x_1(n) W_N^{nk} + e \int_{n=0}^{n-1} x_2(n) W_N^{nk}$
When k is even $e^{\int nk} = 1$
 $x (2k) = \int_{n=0}^{N-1} [x_1(n) + x_2(n)] W_N^{nk}$





IDET USING FET ALGORITHM

$$\begin{aligned} \chi(n) &= \frac{1}{N} \stackrel{N-1}{\underset{k=0}{\overset{K=0}{\underset{k=0}{\atopk=0}{\underset{k=0}{\atopk=0}{\underset{k=0}{\atopk=0}{\underset{k=0}{\atop{k=0}{\atop{k=0}{\atop{k=0}{\atop{k=0}{\atop{k=0}{\atopk=0}{\atop{k=0}{\atopk=0}{\atopk=0}{\atopk=0}{\atopk=0}{\atopk=0}{\atopk=0}{\atop{k=0}{\atop{k=0}{\atop{k=0}{\atop{k=0}{\atop{k=0}{\atop{k=0}{\atop{k=0}{\atop{k=0}{\atop{k=0}{\atopk=0}{\atopk=0}{\atopk=0}{\atopk=0}{\atopk=0}{\atopk=0}{\atopk=0}{\atop{k=0}{\atop{k=0}{\atop{k=0}{\atop{k=0}{\atopk=0}{\atopk=0}{\atopk=0}{\atopk=0}{\atopk=0}{\atopk=0}{\atopk=0}{\atop{k=0}{\atop{k=0}{\atopk=0}{\atopk=0}{\atopk=0}{\atopk=0}{\atopk=0}{\atopk=0}{\atopk=0}{\atop{k=0}{\atopk=0}{$$

Hence IDFT can be found by taking Complexe conjugate and then dividing the sequence by N.

USE OF FFT IN LINEAR FILTERING

Overlap add and Overlap save neethods are used in conjunction with the FFT algorithm for Computing the DFT and the IDFT.

In this, the N-point DFT of him, which is padded by L-A zeros, is denoted as HOA?

This computation is performed once rea the FFT and the resulting N Complex numbers are stoked. DIF-FFT algorithm is used to compute H(K).

Ym(k)=H(k) Xm(k)

Xm(k) -> FFT f x(n) }. Ym(k) in bit reversed order.

The IDFT can be computed by use of an FFT that takes the input is bit-reversed order and produces an output is normal Order. 1) Find the DFT of a sequence $2e(n) = \{1, 2, 3, 4, 4, 3, 2, 1\}$ using DIT algorithm and DIF Algorithm. Solution



DIF ALGORITHM.



OP sequence

$$X(k) = \{20, -5.828, -j2.44, 0, -0.172 - j0.414, 0, -0.172 - j0.414, 0, -0.172 - j0.414, 0, -5.828 - j'2.414\}$$



UNIT-4

STRUCTURES OF IIR (INTRODUCTION.)

Infinite Impube Response (IIR) are recursive type where present output samples depends on the present input, past input and past output samples.

I fulter is one whech rejects the ienwanted frequencies from ip signal and allow the desured frequency to obtain the ofp signal.

Passband :

The range of frequencies that are passed through the fulter are called pass-band.

Ripples:

The linits of tolexance in the magnitude of possiband and stopband are called repples.

Sp -> tolerance in passband

STRUCTURES OF IIR FILTERS

Gystern function of a recursive DDR Puller is given by, H. K

$$H(z) = \frac{\int_{k=0}^{\infty} b_k z^k}{1 + \int_{k=1}^{N} a_k z^{t_1}} - Y^{(0)}$$

DIRECT FORM I STRUCTURES



IRECT FORM REALIZATION OF SYSTEM FUNCTION HICK)

$$H_{2}(x) = \frac{1}{1 + \sum_{k=1}^{N} a_{k} z^{-k}}$$

$$\frac{Y_{2}(x)}{x_{2}(x)} = \frac{1}{1 + \sum_{k=1}^{N} a_{k} z^{-k}}$$

$$Y_{2}(x) \left[1 + \sum_{k=1}^{N} a_{k} x^{-k} \right] = x_{2}(x)$$

$$Y_{2}(x) = -\sum_{k=1}^{N} a_{k} x^{-k} Y_{2}(x) + x_{2}(x)$$

$$Y_{2}(x) = -a_{1} x^{-1} Y_{2}(x) - a_{2} x^{-2} Y_{2}(x) - \dots - a_{N} x^{-N} Y_{2}(x) + x_{2}(x)$$

>



DIRECT FORM-IL STRUCTURE FOR IIR SYSTEM System function of (LTI system) IIR filler is given by $H(Z) = \frac{b_{k} x^{*k}}{1 + b_{k} x^{*k}}$; $y(n) = -\frac{b_{k}}{k} a_{k} y(n-k) + \frac{b_{k}}{k = 0} b_{k} x(n-k) \rightarrow 0$ Let H(X) be written as $H(x) = \frac{Y(x)}{x(x)} = \frac{Y(x)}{W(x)}, \quad \frac{W(x)}{x(x)} = \frac{W(x)}{X(x)}, \quad \frac{Y(x)}{W(x)} = H_1(x), \quad H_2(x)$ $\frac{W(x)}{x(x)} = \frac{1}{1+\sum_{k=1}^{N} a_{k} x^{k}}$ Which given as $W(x) = x(x) - \alpha_1 x^{-1} W(x) - \alpha_2 x^{-2} W(x) - \dots - \alpha_N x^{-N} W(x) \longrightarrow 0$ Y(X) = 5 bx Xth from which $\gamma(x) = (b_0 + b_1 x^{-1} + b_2 x^{-2} + \dots + b_H x^{-H}) W(x)$ +0 By taking Driverse X-transform of equations () and @ $w(n) = x(n) - a_1 bo(n+) - a_2 to(n-2) - \dots - a_N w(n-N) \longrightarrow (5)$ $y(n) = b_{0,+} b_1 w(n-1) + b_2 w(n-2) + \dots + b_N w(n-N) \longrightarrow (5)$ y(n) x(n)w(n) ZT z 6, -ay W(0-1) 1 w(n-1) z1 x-1 62 az w(n-2) w(n-2) w(n-N+1) w(n-N+1) bm-1) z1 -an . by w (n-H) REALIZATION OF EGN (A) REALIZATION OF EGN (S)



MAGINITUDE RESPONSE OF BUTTERWORTH LOWPASS FILTER

Froni the magnitude response plot, the function is . monotonically decreasing, where the maximum response is unity. at 2=0.

The ideal response is shown by the dashed line. The magnitude response approaches the ideal low pars characteri "strics as the order N increases.

For Values 2<2c; [H(j=2)] ×1, for Values 222; [H(j=2)] decreases rapidly. At 2=2c, the curves parses through 0.707, which Corresponds to -3dB point.

In Order to derive transfer function of a stable fulter, substitute S=jer: in ogn O.

$$|H(j,n)|^{2} = H(n^{2}) \qquad |H(j,n)|^{2} = |H(n)| = H(n^{2}) = H(n$$

2

=1,

.: Eq. n. D. well be come

$$H(s) H(s) = \frac{1}{1 + (S_{j})^{2N}}$$

$$= \frac{1}{1 + (E_{j})^{N} s^{2N}}$$

$$H(s) H(s) = \frac{1}{1 + (-s^{2})^{N}}$$

From the above Eqn., H(5) has roots in the LHP and H(5) has the Corresponding roots in the RHP. We can obtain them roots by equating the denominator to zero.

1) Given the specification
$$\alpha_{p} = 1 \text{ dB}$$
; $\alpha_{s} = 30 \text{ dB}$; $\mathcal{D}_{p} = 200 \text{ rad/sec}$
 $\mathcal{D}_{s} = 600 \text{ rad/sec}$. Determine the order of the filter.
Solution
Order of the filter, $N \geq \frac{\log A}{\log k}$
 $A = \frac{A}{2\pi} = \begin{bmatrix} 10^{0.1} \log - 1 \\ 10^{0.1} \log - 1 \end{bmatrix}^{0.5}$
 $A = \frac{A}{2\pi} = \begin{bmatrix} 10^{0.1} \log - 1 \\ 10^{0.1} \log - 1 \end{bmatrix}^{0.5}$
 $A = 63.115$
 $R = \frac{\Omega p}{2\pi} = \frac{200}{600} = \frac{1}{3}$
 $\therefore N \geq \frac{\log 63.115}{\log 3} \geq 3.758$
 $\log 3$
 $\therefore N = 4$
2) Design an analog Butterworth felter that has a - 2 de
Paseband attenuation at a frequency of 20 rad/sec and
at least - u ods stepband attenuation at 30 mad/sec.
Solution:
Given $\alpha_{p} = 2 \text{ dB}$; $\alpha_{s} = 10 \text{ dB}$
 $\Omega_{p} = 20 \text{ rad/sec}$; $\Omega_{s} = 30 \text{ rad/sec}$.
Order of the filter, $N \geq \frac{\log \sqrt{10^{0.05} - 1}}{\log - 2p}$

$$\frac{\log \sqrt{\frac{10^{-1}}{10^{n-1}}}}{\log \frac{36}{252}}$$

$$\geq 3.37$$
Rescueding off N to the next -layboot integer we get
$$N = q.$$
The Normalized Looppers Butters work fully for N24
Can be found written as,
$$H(s) = \frac{1}{(s^2 + 0.765375 + 1)(s^2 + 1.84775 + 1)}$$
(while stephand
$$Le = \frac{-2p}{(10^{n+10}q - 1)^{1/2}N}$$
(while stephand
$$-2e = \frac{-2p}{(10^{n+10}q - 1)^{1/2}N}$$
(while stephand
$$-2e = \frac{-2p}{(10^{n+10}q - 1)^{1/2}N}$$
Here step band operification at 2s is exceeded, so
$$-2e = \frac{-2p}{(10^{n+10}q - 1)^{1/2}N} = 21.3868$$
The Alassfer function for $P_{z} = d1.3868$ can be
obtained by Substituting
$$S \rightarrow \frac{2}{-2e} = \frac{1}{21.3868} + 1(s)$$

$$H(s) = \frac{1}{(s + 10^{1/2} - 1)^{1/2}N} + (s + 10^{1/2} - 1)^{1/2}N$$

$$Res = 1$$

$$H(s) = \frac{1}{(s + 10^{1/2} - 1)^{1/2}N} + 1(s + 10^{1/2} - 1)^{1/2}N$$

$$Res = 1$$

$$Res = 1$$

$$H(s) = \frac{1}{(s + 10^{1/2} - 1)^{1/2}N} + 1(s + 10^{1/2} - 1)^{1/2}N$$

$$Res = 1$$

STEPS TO DESIGN A DIGITAL FILTER USING IMPULSE INVARIANCE METHOD. 14 For the green specifications, find Ha(S). 24; Solect the Sampling rate of the digital feller, 7 Seconds per Sample. 34 Eseptess the analog filter theoretic functions as the sum of Single pole filters Ha(S) = $\sum_{k=1}^{N} \frac{Ck}{S-P_k}$

17 Compute the z-triansform of the digital filter

$$H(z) = \sum_{k=1}^{N} \frac{C_{k}}{1 - e^{T_{k}T_{z}-1}}$$

For frigh Campling rate
$$H(z) = \sum_{k=1}^{N} \frac{T_{k}C_{k}}{1 - e^{T_{k}T_{z}-1}}$$

For the analog transfer function
$$H(s) = \frac{2}{(S+A)^{(S+A)}}$$
 determine
 $H(x)$ using impulse invariance method. decrement $T=1$ secs
solution:
 $G(treen H(s)) = \frac{2}{(S+1)^{(S+A)}}$
Using portial fraction
 $H(s) = \frac{A}{S+1} + \frac{B}{S+2}$
 $a = A(S+A) + B(S+1)$
Sub $S=-1$; $a = cA(1)$
 $cA = 2$
 $Sub S=-2$; $b = B(1)$
 $B=-2$
 $H(S) = \frac{2}{S-1} - \frac{2}{S-(-2)}$
Using impulse invariance technique, if
 $H(s) = \frac{A}{S-R} + \frac{CR}{S-R}$ then $H(x) = \frac{A}{K+1} + \frac{CR}{1-e^{2R}x^{-1}}$.
(a) $(s+R_1)$ is transformed site $1-e^{RT}x^{-1}$.
There are have poles $P_1 = H + 2 P_2 = -3$. SD
 $H(x) = \frac{2}{1-e^{-T}x^{-1}} - \frac{2}{1-e^{2R}x^{-1}}$
 $For T = 1$ sec
 $H(x) = \frac{2}{1-e^{-T}x^{-1}} - \frac{2}{1-e^{2R}x^{-1}}$
 $H(x) = \frac{2}{1-e^{-T}x^{-1}} - \frac{2}{1-e^{2R}x^{-1}}$
THE WARPING EFFECT Let \mathcal{D} and \mathcal{W} represent. the frequency: Variables in the analog folter and the derived digital folter respectively. $\mathcal{D} = \frac{2}{7} \tan \frac{10}{2}$. For small value of \mathcal{W} $\mathcal{D} = \mathcal{D} = \frac{2}{7}$. $\mathcal{W} = \mathcal{W} = \frac{10}{7}$ $\mathcal{W} = \mathcal{D} = \frac{10}{7}$ $\mathcal{W} = \mathcal{D} = 0$.



sy Soleet the Soupleig cab of the degrial fulles;
Call it is second per souple
As Substitute
$$s_{2-x} = \frac{1-x^{2}}{1+x^{2}}$$
 into the transfer function
found in step 2.

Apply bettinear transformation to $H(s) = \frac{2}{(s+1)(s+2)}$ with $T \ge 1$ sec
and Find $H(x)$.
Solution
 $f(x) = \frac{2}{(s+1)(s+2)}$
Substitute $s = \frac{2}{r} \left[\frac{1-x^{2}}{1+x^{2}} \right]$ is $H(s)$ to get $H(x)$
 $H(x) \ge H(s) / s = \frac{2}{r} \left[\frac{1-x^{2}}{1+x^{2}} \right]$
 $H(x) = \frac{2}{(s+1)(s+2)} / s = \frac{2}{r} \left[\frac{1-x^{2}}{1+x^{2}} \right]$
 $H(x) = \frac{2}{(s+1)(s+2)} / s = \frac{2}{r} \left[\frac{1-x^{2}}{1+x^{2}} \right]$
 $T \ge 1$ sec.
 $H(x) = \frac{2}{\left\{ \frac{2}{n} \left(\frac{1-x^{2}}{1+x^{2}} \right\} + 1 \right\} \left\{ \frac{2}{n} \left(\frac{1-x^{2}}{1+x^{2}} \right\} + 2 \right\}$
 $= \frac{2(1+x^{2})^{2-1}}{(2-3x^{2}+1+x^{2})(2-3x^{2}+2+3x^{2})}$
 $= \frac{2(1+x^{2})^{2}}{(2-3x^{2}+1+x^{2})(2-3x^{2}+2+3x^{2})}$
 $= \frac{2(1+x^{2})^{2}}{(2-3x^{2}+1+x^{2})(2-3x^{2}+2+3x^{2})}$
 $= \frac{2(1+x^{2})^{2}}{(2-3x^{2}+1+x^{2})(2-3x^{2}+2+3x^{2})}$
 $= \frac{2(1+x^{2})^{2}}{(2-3x^{2}+1+x^{2})(2-3x^{2}+2+3x^{2})}$
 $= \frac{2(1+x^{2})^{2}}{(1-5x^{2}-3x^{2})}$ or
 $H(x) = \frac{-(16k(1+x^{2}))^{2}}{(1-5x^{2}-3x^{2})}$

DESIGN OF HIGHPASS, BANDPASS AND EANDSTOP FILTERS
To find the transfer function of highpan,
bandpass and band stop filters of any type fast
find the transfer function of nernalized 2 PF. (Mass
and then we surfable transferration is used.

For the given specifications
$$\alpha_p = 3dB$$
; $\alpha_s = 16dE$, $\Lambda p = 1000$ rad/sec
and $\Lambda_s = 500$ rad/sec design a highpan filter.
Solution:
First we design a normalized LPF and theo suitable
transferration is used to get the transfer function of
a highpars filter.
For LPF
 $\Lambda_c = \Lambda p = 500$ rad/sec
 $-\Omega_s = 1000$ rad/sec
 $\Omega_s = 0.5$
 $\Omega = \frac{\log S}{\Lambda s} = 2.060$
 $N = 3$ CEY approximity Next higher integes
H(00 for $\Omega_s = 1$ rad/sec and N=2 B
 $H(00 = \frac{1}{(B+1)(S^2+S+1)})$

To get highpans filter having cut-off frequency

$$-\Omega_{c} = \Omega_{p} = 1000 \text{ rad/sec}.$$
Sub $s \rightarrow \frac{1000}{s}$

$$H_{a}(s) = H(s) / \frac{1000}{s} = \frac{1}{(s+1)(s^{2}+s+1)} / \frac{1000}{s} = \frac{1}{(s+1)(s^{2}+s+1)} / \frac{1000}{s} = \frac{1}{(\frac{1000}{s}+1)\left[\frac{(1000)^{2}}{s^{2}} + \frac{1000}{s} + \frac{1}{s}\right]}$$

$$= \frac{10}{(\frac{1000+s}{s})\left(\frac{(1000)^{2}+1000s+s^{2}}{s^{2}}\right)}$$

$$H_{a}(s) = \frac{s^{3}}{(s+1000)(s^{2}+1000s+(1000s^{2}+1000s^{2}))}$$

ANALOG LOWPASS CHEBYSHEV FILTERS These are two types of chebysher fultus. Type I chebysher falters are all-pole falters that exhibit equiripple behavious in the passband and a nonotonic characteristrics in the stopband. Type I chebysher fulter contains both poles and zeros and exhibits a monotonic behaviour in the passband and an equilipple behaviour is the stop band. [40.2)] 1 HOLDI 146 J1-172 -2p -2s TYPE I TYPE I CHARACTERISTICS OF CHEBYSHEV FILTERS 1HCJ-0 1400 VH4 Xs JHS Ks J1-172 2p-Sz 2p-52c -LOWPASS CHEBYSHEV FILTER MAGNITUDE RESPONSE

STEPS TO DESIGN AN ANALOG CHEBYSHEV LOWPASS FILT
1. From the given specifications find the order of
the filler Nussing the formula

$$N \ge (\cos h^{-1} \int \frac{10^{\circ 108} - 1}{10^{\circ 108} - 1}) / (\cosh^{-1} (28))$$

Where $\cosh^{-1} x = \ln (x + \sqrt{x^2 - 1}) = \frac{\log (x + \sqrt{x^2 - 1})}{\log 2}$
2. Round off it to the next higher integer.
3 For even values of N the oscillatory curves
starts from $\int \frac{1}{\sqrt{1+5t^2}}$
For odd values of N the oscillatory curves starts
from 1.
3. Using the following formulas find the values
of a and b, which are minor and waits on the

a=
$$\Omega_p\left[\frac{\mu^{N}}{2}-\frac{\mu^{-N}}{2}\right]$$
; $b=\Omega_p\left[\frac{\mu^{N}+\mu^{-N}}{2}\right]$

$$\mu = \varepsilon_1^{-1} + \sqrt{\varepsilon_1^{-2}}$$

the

-2p= Passband frequency xp = Marenum allowable attenuation in the passbance [::For normalized chebysher fultu sp=1 xad/see 4. Calculate the potes of chebysher fultu which lie on an ellipse by using the formula

SK= a cos of tibsin of K=1,2,... N Where $\phi_{k} = \frac{\pi}{2} + \left(\frac{2\kappa - 1}{2N}\right) \pi$, $\kappa = 1, 2, ... N$ 5. Find the denominator polynomial of the bransfer function using the above potes 1'2> $(S_k - a \cos \phi_k)^2 + (b \sin \phi_k)^2$ 6. The numerator of the biarsfer functions depends or the value of N. a> For N odd Substitute s=0 in the denomination polynomial and find the value. This value is equal to the numerator of the transfer function. by for N even substitute s=0 is the denominato polynomial and divede the result by VI+52. Thus Value is equal to the numerator. 1. Given the specifications xp=3 dB; xs=16 dB; fp=1 KHz and fs=2 KHZ. Determine the order of the filter using chebysher approximations. Find H(s). Solution From the given data we can find Ip=2xx1000Hz=2000x rad/sec -25 = 2 x x 2000Hz = 4000 x rad/sec and xp=3 dB ; xs=16 dB i' order of the fully, N $N \ge \frac{\cosh^{-1} \sqrt{10^{0.1} \alpha_{s-1}}}{\sqrt{10^{0.1} \alpha_{P-1}}} \ge \frac{\cosh^{-1} \sqrt{10^{0.3} - 1}}{\cosh^{-1} (\frac{4000 \pi}{2000 \pi})}$

$$\sum_{k=1}^{\infty} \frac{\cosh^{-1}(6\cdot 2q_{k})}{\cosh^{-1}(2\cdot)}$$

$$\sum_{k=1}^{\infty} \frac{\ln\left[6\cdot 2q_{k}q_{k}\sqrt{6\cdot 2q_{k}q_{k}^{2}-1}\right]}{\ln\left[2+\sqrt{2^{2}+1}\right]}$$

$$\sum_{k=1}^{\infty} \frac{3\cdot 6\cdot 332}{1\cdot 8\cdot 9\cdot q}$$

$$N \ge 1.91$$

$$\sum_{k=1}^{\infty} \frac{3\cdot 6\cdot 332}{1\cdot 8\cdot 9\cdot q}$$

$$N \ge 1.91$$

$$\sum_{k=1}^{\infty} \frac{3\cdot 6\cdot 332}{1\cdot 8\cdot 9\cdot q}$$

$$N \ge 1.91$$

$$\sum_{k=1}^{\infty} \frac{1}{1\cdot 8\cdot q}$$

$$S_{1} = \alpha \cos \phi_{1} + jbsin \phi_{1} = -643 \cdot 4b\pi + j1554\pi$$

$$S_{2} = \alpha \cos \phi_{2} + jbsin \phi_{2} = -643 \cdot 4b\pi - j1554\pi$$

$$S_{2} = \alpha \cos \phi_{2} + jbsin \phi_{2} = -643 \cdot 4b\pi - j1554\pi$$

$$S_{1} \cdot The penonuinator of H(s) = (S+643 \cdot 4b\pi)^{2} + (1554\pi)^{2}$$

$$6Y \cdot The Numerator of H(s) = -\frac{(643 \cdot 4b\pi)^{2} + (1554\pi)^{2}}{\sqrt{1+g_{1}^{2}}} \quad [: N beaun]$$

$$= (1414 \cdot 38)^{2}\pi^{2}$$

$$= (1414 \cdot 38)^{2}\pi^{2}$$

$$H(s) = \frac{(1414 \cdot 38)^{2}\pi^{2}}{s^{2} + 1287\pi s + (1682)^{2}\pi^{2}}$$

UNIT-5

DIGITAL SIGNAL PROCESSORS

ARCHITECTURE

The '54x DSPs use an advanced, modified Harvard architecture that maximizes processing power by maintaining one program memory bus and three data memory buses. These processors also provide an arithmetic logic unit (ALU) that has a high degree of parallelism, application-specific hardware logic, on-chip memory, and additional on-chip peripherals. These DSP families also provide a highly specialized instruction set, which is the basis of the operational flexibility and speed of these DSPs. Separate program and data spaces allow simultaneous access to program instructions and data, providing the high degree of parallelism. Two reads and one write operation can be performed in a single cycle. Instructions with parallel store and application-specific instructions can fully utilize this architecture. In addition, data can be transferred between data and program spaces. Such parallelism supports a powerful set of arithmetic, logic, and bit-manipulation operations that can all be performed in a single machine cycle. Also included are the control mechanisms to manage interrupts, repeated operations, and function calls. Figure 1 shows the functional block diagram that shows the principal blocks and bus structure in the '54x devices

TMS320C5'416 FAMILY – FUNCTIONAL OVERVIEW Central Processing Unit (CPU)

The CPU of the '54x devices contains:

□ A 40-bit arithmetic logic unit (ALU)

- □ Two 40-bit accumulators
- □ A barrel shifter
- □ A 17 × 17-bit multiplier/adder

Arithmetic Logic Unit (ALU)

The '54x devices perform 2s-complement arithmetic using a 40-bit ALU and two 40-bit accumulators (ACCA and ACCB). The ALU also can perform Boolean operations. The ALU can function as two 16-bit ALUs and perform two 16-bit operations simultaneously when the C16 bit in status register 1 (ST1) is set.



Accumulators

The accumulators, ACCA and ACCB, store the output from the ALU or the multiplier / adder block; the accumulators can also provide a second input to the ALU or the multiplier / adder. The bits in each accumulator is grouped as follows: □ Guard bits (bits 32–39)

- \Box A high-order word (bits 16–31)
- \Box A low-order word (bits 0–15)

Instructions are provided for storing the guard bits, the high-order and the low-order accumulator words in data memory, and for manipulating 32-bit accumulator words in or out of data memory. Also, any of the accumulators can be used as temporary storage for the other.

Barrel Shifter

The '54x's barrel shifter has a 40-bit input connected to the accumulator or data memory (CB, DB) and a 40-bit output connected to the ALU or data memory (EB). The barrel shifter produces a left shift of 0 to 31 bits and a right shift of 0 to 16 bits on the input data. The shift requirements are defined in the shift-count field (ASM) of ST1 or defined in the temporary register (TREG), which is designated as a shift-count register. This shifter and the exponent detector normalize the values in an accumulator in a single cycle. The least significant bits (LSBs) of the output are filled with 0s and the most significant bits (MSBs) can be either zero-filled or sign-extended, depending on the state of the sign-extended mode bit (SXM) of ST1. Additional shift capabilities enable the processor to perform numerical scaling, bit extraction, extended arithmetic, and overflow prevention operations.

Multiplier/Adder

The multiplier / adder performs 17 x 17-bit 2s-complement multiplication with a 40-bit accumulation in a single instruction cycle. The multiplier / adder block consists of several elements; a multiplier, adder, signed/unsigned input control, fractional control, a zero detector, a rounder (2s-complement), overflow/saturation logic, and TREG. The multiplier has two inputs: one input is selected from the TREG, a data-memory operand, or an accumulator; the other is selected from the program memory, the data memory, an accumulator, or an immediate value. The fast on-chip multiplier allows the '54x to perform operations such as convolution, correlation, and filtering efficiently. In addition, the multiplier and ALU together execute multiply/accumulate (MAC) computations and ALU operations in parallel in a single instruction cycle. This function is used in determining the Euclid distance, and in implementing symmetrical and least mean square (LMS) filters, which are required for complex DSP algorithms.

Compare, Select, and Store Unit (CSSU)

The compare, select, and store unit (CSSU) performs maximum comparisons between the accumulator's high and low words, allows the test/control (TC) flag bit of status register 0 (ST0) and the transition (TRN) register to keep their transition histories, and selects the larger word in the accumulator to be stored in data memory. The CSSU also accelerates Viterbi-type butterfly computation with optimized on-chip hardware.

Program Control

Program control is provided by several hardware and software mechanisms:

□ The program controller decodes instructions, manages the pipeline, stores the status of operations, and decodes conditional operations. Some of the hardware elements included in the program controller are the program counter, the status and control register, the stack, and the address-generation logic.

□ Some of the software mechanisms used for program control include branches, calls, conditional instructions, a repeat instruction, reset, and interrupts.

□ The '54x supports both the use of hardware and software interrupts for program control. Interrupt service routines are vectored through a relocatable interrupt vector table. Interrupts can be globally enabled/disabled and can be individually masked through the interrupt mask register (IMR). Pending interrupts are indicated in the

interrupt flag register (IFR). For detailed information on the structure of the interrupt vector table, the IMR and the IFR, see the device-specific data sheets.

Status Registers (ST0, ST1)

The status registers, ST0 and ST1, contain the status of the various conditions and modes for the '54x devices. ST0 contains the flags (OV, C, and TC) produced arithmetic operations and bit manipulations in addition to the data page pointer (DP) and the auxiliary register pointer (ARP) fields. ST1 contains the various modes and instructions that the processor operates on and executes.

Auxiliary Registers (AR0–AR7)

The eight 16-bit auxiliary registers (AR0–AR7) can be accessed by the central arithmetic logic unit (CALU) and modified by the auxiliary register arithmetic units (ARAUs). The primary function of the auxiliary registers is generating 16-bit addresses for data space. However, these registers also can act as general-purpose registers or counters.

Temporary Register (TREG)

The TREG is used to hold one of the multiplicands for multiply and multiply/accumulate instructions. It can hold a dynamic (execution-time programmable) shift count for instructions with a shift operation such as ADD, LD, and SUB. It also can hold a dynamic bit address for the BITT instruction. The EXP instruction stores the exponent value computed into the TREG, while the NORM instruction uses the TREG value to normalize the number. For ACS operation of Viterbi decoding, TREG holds branch metrics used by the DADST and DSADT instructions.

Transition Register (TRN)

The TRN is a 16-bit register that is used to hold the transition decision for the path to new metrics to perform the Viterbi algorithm. The CMPS (compare, select, max, and store) instruction updates the contents of the TRN based on the comparison between the accumulator high word and the accumulator low word.

Stack-Pointer Register (SP)

The SP is a 16-bit register that contains the address at the top of the system stack. The SP always points to the last element pushed onto the stack. The stack is manipulated by interrupts, traps, calls, returns, and the PUSHD, PSHM, POPD, and POPM instructions. Pushes and pops of the stack predecrement and postincrement, respectively, all 16 bits of the SP.

Circular-Buffer-Size Register (BK)

The 16-bit BK is used by the ARAUs in circular addressing to specify the data block size. **Block-Repeat Registers (BRC, RSA, REA)**

The block-repeat counter (BRC) is a 16-bit register used to specify the number of times a block of code is to be repeated when performing a block repeat. The block-repeat start address (RSA) is a 16-bit register containing the starting address of the block of program memory to be repeated when operating in the repeat mode. The 16-bit block-repeat end address (REA) contains the ending address if the block of program memory is to be repeated when operating in the repeat mode.

Interrupt Registers (IMR, IFR)

The interrupt-mask register (IMR) is used to mask off specific interrupts individually at required times. The interrupt-flag register (IFR) indicates the current status of the interrupts.

Processor-Mode Status Register (PMST)

The processor-mode status register (PMST) controls memory configurations of the '54x devices.

Power-Down Modes

There are three power-down modes, activated by the IDLE1, IDLE2, and IDLE3 instructions. In these modes, the '54x devices enter a dormant state and dissipate considerably less power than in normal operation. The IDLE1 instruction is used to shut down the CPU. The IDLE2 instruction is used to shut down the CPU and on-chip peripherals. The IDLE3 instruction is used to shut down the '54x processor completely. This instruction stops the PLL circuitry as well as the CPU and peripherals.

BUS STRUCTURE

The '54x device architecture is built around eight major 16-bit buses: □ One program-read bus (PB) which carries the instruction code and immediate operands from program memory

□ Two data-read buses (CB, DB) and one data-write bus (EB), which interconnect to various elements, such as the CPU, data-address generation logic (DAGEN), program-address generation logic (PAGEN), on-chip peripherals, and data memory

□ The CB and DB carry the operands read from data memory.

• The EB carries the data to be written to memory.

□ Four address buses (PAB, CAB, DAB, and EAB), which carry the addresses needed for instruction execution The '54x devices have the capability to generate up to two datamemory addresses per cycle, which are stored into two auxiliary register arithmetic units (ARAU0 and ARAU1).

The PB can carry data operands stored in program space (for instance, a coefficient table) to the multiplier for multiply /accumulate operations or to a destination in data space for the data-move instruction. This capability allows implementation of single-cycle three-operand instructions such as FIRS. The '54x devices also have an on-chip bidirectional bus for accessing on-chip peripherals; this bus is connected to DB and EB through the bus exchanger in the CPU interface. Accesses using this bus can require more than two cycles for reads and writes depending on the peripheral's structure. The '54x devices can have bus holders connected to the data bus and the HPI data bus. Bus holders ensure that the data bus does not float. When bus holders are enabled, the data bus maintains its previous level. Setting bit 1 of the bank-switching control register (BSCR) enables bus holders and clearing bit 1 disables the bus holders. A reset automatically disables the bus holders. The bus holders ensure that the address bus. The bus holders ensure that the address bus holders ensure that the address bus. The bus holders ensure that the address bus does not float when in high impedance. For these devices, the bus holders are always enabled.

MEMORY

The minimum memory address range for the '54x devices is 192K words — composed of 64K words in program space, 64K words in data space, and 64K words in I/O space. Selected devices also provide extended program memory space of up to 8M words. The program memory space contains the instructions to be executed as well as tables used in execution. The data memory space stores data used by the instructions. The I/O memory space interfaces to external memory-mapped peripherals and can also serve as extra data storage space. The '54x DSPs provide both on-chip RAM and ROM to improve system performance and integration.

On-Chip ROM

The '54x devices include on-chip maskable ROM that can be mapped into program memory or data memory depending on the device. On-chip ROM is mapped into program space by the microprocessor/microcontroller (MP/MC) mode control pin. On-

chip ROM that can be mapped into data space is controlled by the DROM bit in the processor mode status register (PMST). This allows an instruction to use data stored in the ROM as an operand. Customers can arrange to have the ROM of the '54x programmed with contents unique to any particular application.

Bootloader

A bootloader is available in the standard '54x on-chip ROM. This bootloader can be used to transfer user code from an external source to anywhere in the program memory at power up automatically. If the MP/MC pin of the device is sampled low during a hardware reset, execution begins at location FF80h of the on-chip ROM. This location contains a branch instruction to the start of the bootloader program.

On-Chip Dual-Access RAM (DARAM)

Dual-access RAM blocks can be accessed twice per machine cycle. This memory is intended primarily to store data values; however, it can be used to store program as well. At reset, the DARAM is mapped into data memory space. DARAM can be mapped into program/data memory space by setting

the OVLY bit in the PMST register.

On-Chip Single-Access RAM (SARAM)

Each of the SARAM blocks is a single-access memory. This memory is intended primarily to store data values; however, it can be used to store program as well. SARAM can be mapped into program/data memory space by setting the OVLY bit in the PMST register.

On-Chip Two-Way Shared RAM

Select 54x devices with multiple CPU cores include two-way shared RAM blocks that allow simultaneous program space access from two CPU cores. Each CPU can perform a single access with zero-states to any location in the two-way shared RAM during each clock cycle. This shared RAM is most efficiently used when the two CPUs are executingidentical programs. In this case, the amount of program memory required for the application is effectively reduced by 50% since both CPUs can execute from the same RAM.

On-Chip Memory Security

A security feature is included on 54x devices to prevent the on-chip memory contents from being extracted by a user. This feature is enabled during the manufacturing process and is ONLY available to customers that order custom ROM programming. Consequently, memory security cannot be enabled/disabled by the user.

Program Memory

The standard external program memory space on the '54x devices addresses up to 64K 16-bit words. Software can configure their memory cells to reside inside or outside of the program address map. When the cells are mapped into program space, the device automatically accesses them when their addresses are within bounds. When the program-address generation (PAGEN) logic generates an address outside its bounds, the device automatically generates an external access.

PIPELINING

Instruction pipelining is a technique that implements a form of parallelism called instruction-level parallelism within a single processor. It therefore allows faster CPU throughput (the number of instructions that can be executed in a unit of time) than would otherwise be possible at a given clock rate. The basic instruction cycle is broken up into a series called a pipeline. Rather than processing each instruction sequentially (finishing one instruction before starting the next), each instruction is split up into a sequence of steps so different steps can be executed in parallel and instructions can be processed concurrently (starting one instruction before finishing the previous one). identical programs. In this case, the amount of program memory required for the application is effectively reduced by 50% since both CPUs can execute from the same RAM.

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Indirect Addressing

• Data space is accessed by address present in an auxiliary register. • 54xx have 8, 16 bit auxiliary register (AR0 – AR 7). Two auxiliary register arithmetic units (ARAU0 & ARAU1) • Used to access memory location in fixed step size. AR0 register is used for indexed and bit reverse addressing modes. • For single – operand addressing MOD type of indirect addressing ARF AR used for addressing

• ARP depends on (CMPT) bit in ST1 CMPT = 0, Standard mode, ARP set to zero CMPT = 1, Compatibility mode, Particularly AR selected by ARP



Table 1 Indiect addressing option with single data memory operand ; Circular addressing\

Operand syntax	Function
*ARx	Addr = ARx;
*ARx -	Addr = ARx; $ARx = ARx - 1$
*ARx +	Addr = ARx; ARx = ARx +1
*+ARx	Addr = ARx+1; ARx = ARx+1
*ARx - 0B	Addr = ARx ; $ARx = B(ARx - AR0)$
*ARx - 0	Addr = Arx; $ARx = ARx - AR0$
*ARx + 0	Addr = Arx; $ARx = ARx + AR0$
*ARx + 0B	Addr = ARx; $ARx = B(ARx + AR0)$
*ARx - %	Addr = ARx ; ARx = circ(ARx - 1)

Bit-Reversed Addressing:

• Used for FFT algorithms. • AR0 specifies one half of the size of the FFT. • The value of AR0 = 2N-1: N = integer FFT size = 2N • AR0 + AR (selected register) = bit reverse addressing. • The carry bit propagating from left to right.

Dual-Operand Addressing

Dual data-memory operand addressing is used for instruction that simultaneously perform two reads (32-bit read) or a single read (16-bit read) and a parallel store (16-bit store) indicated by two vertical bars, II. These instructions access operands using indirect addressing mode.

If in an instruction with a parallel store the source operand the destination operand point to the same location, the source is read before writing to the destination. Only 2 bits are available in the instruction code for selecting each auxiliary register in this mode. Thus, just four of the auxiliary registers, AR2-AR5, can be used, The ARAUs together with these registers, provide capability to access two operands in a single cycle. Figure 6 shows how an address is generated using dual data-memory operand addressing.



Memory-Mapped Register Addressing

• Used to modify the memory-mapped registers without affecting the current datapage pointer (DP) or stack-pointer (SP) – Overhead for writing to a register is minimal – Works for direct and indirect addressing - Scratch -pad RAM located on data PAGE0 can be modified • STM #x, DIRECT • STM #tbl, AR1

Figure



Stack Addressing • Used to automatically store the program counter during interrupts and subroutines. • Can be used to store additional items of context or to pass data values. • Uses a 16-bit memory-mapped register, the stack pointer (SP). • PSHD X2